

# Double Fibrations



Since it is not obvious what the codomain  
of an indexing functor

$$\mathcal{A} \longrightarrow ?? \quad (\mathcal{A} \text{ dbl cat}^{\mathcal{F}})$$

should be, we start with double fibrations.



# Double Fibrations

First Idea: a double category is a pseudo cat<sup>d</sup> in Cat and a double fibration is a pseudo cat<sup>d</sup> in Fib.

This would amount to:

$$\begin{array}{ccccc}
 \underline{E}_1 \times_{\underline{E}_0} \underline{E}_1 & \xrightarrow{\otimes^T} & \underline{E}_1 & \begin{array}{c} \xrightarrow{s^T} \\ \xleftarrow{y^T} \\ \xrightarrow{t^T} \end{array} & \underline{E}_0 \\
 \downarrow \begin{array}{c} P_1 \times P_1 \\ P_0 \end{array} & & \downarrow P_1 & & \downarrow P_0 \\
 \underline{B}_1 \times_{\underline{B}_0} \underline{B}_1 & \xrightarrow{\otimes^L} & \underline{B}_1 & \begin{array}{c} \xrightarrow{s^L} \\ \xleftarrow{y^L} \\ \xrightarrow{t^L} \end{array} & \underline{B}_0
 \end{array}$$

... strict dbl functor with extra properties?



## Double fibrations

Problem : Fib doesn't have all 2-pullbacks required for this.

Observation:

- We require the same fibrational strictness for  $s$  and  $t$  that we require for  $y$  and  $\otimes$ .

Solution: We will require that  $s$  and  $t$  are in  $\mathcal{CFib}$ .  
(i.e., they preserve cleavages)



# Double Fibrations

Definition: A double fibration is a (strict) double functor  $P: \mathbb{E} \rightarrow \mathbb{B}$  between pseudo double categories:

$$\begin{array}{ccc}
 \underline{E}_1 \times_{\underline{E}_0} \underline{E}_1 & \xrightarrow{\otimes_E} & \underline{E}_1 \begin{array}{c} \xrightarrow{s_E} \\ \xleftarrow{t_E} \end{array} \underline{E}_0 \\
 \downarrow \begin{array}{c} P_1 \\ P_0 \end{array} & & \downarrow P_i \\
 \underline{B}_1 \times_{\underline{B}_0} \underline{B}_1 & \xrightarrow{\otimes_B} & \underline{B}_1 \begin{array}{c} \xrightarrow{s_B} \\ \xleftarrow{t_B} \end{array} \underline{B}_0
 \end{array} \quad \text{s.t.}$$

- $P_0$  and  $P_1$  are fibrations with cleavages
- $s_E$  and  $t_E$  are cleavage preserving
- $y_E$  and  $\otimes_E$  are cartesian-morphism preserving.



## Examples

1.  $\text{Im} : \mathcal{S}\text{pan}(\underline{\text{Set}}) \rightarrow \text{Rel}$   
 $(A \xleftarrow{s} S \xrightarrow{t} B) \mapsto (\text{Im}(s,t) \subseteq A \times B)$   
 is a double opfibration.

2. When  $\underline{E}_0 = \underline{B}_0 = 1$ , we get:

$$\begin{array}{ccccc}
 \underline{E}_1 \times_{\underline{E}_0} \underline{E}_1 & \xrightarrow{\otimes_E} & \underline{E}_1 & \xleftarrow{\eta} & 1 \\
 \downarrow \text{P}_1 \times \text{P}_1 & & \downarrow \text{P}_1 & & \downarrow \text{id} \\
 \underline{B}_1 \times_{\underline{B}_0} \underline{B}_1 & \xrightarrow{\otimes_B} & \underline{B}_1 & \xleftarrow{\eta} & 1
 \end{array}$$

a monoidal fibration!



3. The Grothendieck constructions given in David Jaz Myers' work are also double fibrations.

## Examples

4. For any 2-functor  $P: \underline{E} \rightarrow \underline{B}$ ,  $P$  is a 2-fibration as in Buckley's work if and only if  $\mathbb{Q}(P): \mathbb{Q}(E) \rightarrow \mathbb{Q}(B)$  is a double fibration.

5. if  $P_0$  and  $P_1$  are discrete fibrations, we recover discrete double fibrations.

6. For  $\mathbb{D}$  a double cat<sup>y</sup>, let  $\mathbb{D}^z = \begin{pmatrix} \mathbb{D}_1^z \\ \Downarrow \\ \mathbb{D}_0^z \end{pmatrix}$

dom:  $\mathbb{D}^z \rightarrow \mathbb{D}$

is a double fibration



7. The codomain fibration extends to a double codomain fibration  $\text{cod}: \mathbb{D}^2 \rightarrow \mathbb{D}$  if:

- $\mathbb{D}_1$  and  $\mathbb{D}_0$  have chosen finite limits
- these limits are preserved on the nose by  $s$  and  $t$
- and up to iso by  $y$  and  $\otimes$ .





8. for  $\mathcal{C}$  a small cat<sup>d</sup>,  $\text{Fam}(\mathcal{C})$  has:

obj:  $f: I \rightarrow \mathcal{C}$ , or  $(I, \{C_i\}_{i \in I})$

arrows:  $(h, \alpha): f \rightarrow g$   $I \xrightarrow{h} J$   
 $h: I \rightarrow J, \alpha: f \Rightarrow g \circ h$   $\begin{array}{ccc} & \xrightarrow{\alpha} & \\ f & \searrow & \swarrow g \\ & \circlearrowright & \end{array}$

or:  $(I, \{C_i\}_{i \in I})$

↓

$(h: I \rightarrow J, \{\alpha_i: C_i \rightarrow C_{h(i)}\}_{i \in I})$

↓

$(J, \{C_j\}_{j \in J})$



Proarrows:  $f \xrightarrow{(S, d_0, d_1, \theta)}$

for natural transformations

$$\begin{array}{ccc} S & \xrightarrow{d_1} & K \\ d_0 \downarrow \theta & \Rightarrow & \downarrow p \\ I & \xrightarrow{f} & e \end{array}$$

for a span of functions

$$I \xleftarrow{d_0} S \xrightarrow{d_1} K$$

or:

$$(I, \{C_i\}_{i \in I}) \xrightarrow{((d_0, S, d_1), \theta)} (K, \{C_k\})$$

for a span of functions  $I \xleftarrow{d_0} S \xrightarrow{d_1} K$

+ a family of arrows

$$\theta_S : C_{d_0(s)} \longrightarrow C_{d_1(s)}$$



cells: a cell from

$$\begin{array}{ccc}
 S & \xrightarrow{d_1} & K \\
 d_0 \downarrow \cong \Downarrow & & \downarrow \cong \\
 I & \xrightarrow{f} & e
 \end{array}
 \quad + \quad
 \begin{array}{ccc}
 T & \xrightarrow{d'_1} & L \\
 d'_0 \downarrow \cong \Downarrow & & \downarrow \cong \\
 J & \xrightarrow{g} & e
 \end{array}$$

is given by a morphism of spans:

$$\begin{array}{ccccc}
 I & \xleftarrow{d_0} & S & \xrightarrow{d_1} & K \\
 h \downarrow & & \downarrow m & & \downarrow r \\
 J & \xleftarrow{d'_0} & T & \xrightarrow{d'_1} & L
 \end{array}$$

with 2-cells:

$$\begin{array}{ccc}
 & I & \\
 f \swarrow & & \searrow h \\
 e & \xrightarrow{\alpha} & \\
 g \swarrow & & \searrow J
 \end{array}$$

$$\begin{array}{ccc}
 & K & \\
 p \swarrow & & \searrow r \\
 e & \xrightarrow{\beta} & \\
 q \swarrow & & \searrow L
 \end{array}$$

st:

$$(\beta * d_1) \theta =$$

$$(\delta * g) (\alpha * d_0)$$



or: a family of cells:

$$\begin{array}{ccc} (I, \{c_i\}_{i \in I}) & \xrightarrow{((d_0, s, d_1), \theta)} & (K, \{c_k\}_{k \in K}) \\ (h, (\alpha_i)) \downarrow & m & \downarrow (r, (\beta_k)) \\ (J, \{c_j\}_{j \in J}) & \xrightarrow{((d'_0, T, d'_1), \theta')} & (L, \{c_l\}_{l \in L}) \end{array}$$

where  $m: S \rightarrow T$  fits in

$$\begin{array}{ccccc} I & \xleftarrow{d_0} & S & \xrightarrow{d_1} & K \\ h \downarrow & = & \downarrow m & \neq & \downarrow r \\ J & \xleftarrow{d'_0} & T & \xrightarrow{d'_1} & L \end{array}$$

and we require that for each  $s \in S$ :



$$\begin{array}{ccc}
 C_{d_0}(s) & \xrightarrow{\Theta_s} & C_{d_1}(s) \\
 \alpha_{d_1(s)} \downarrow & & \downarrow \beta_{d_1(s)} \\
 C_{h(d_0, s)} & \xrightarrow{\Theta'_{m(s)}} & C_{r(d_1, s)}
 \end{array}$$

this square is well defined because

$$d_0'(m(s)) = h(d_0, s)$$

$$\text{and } d_1'(m(s)) = r(d_1, s)$$

we only need to require that

$$\beta_{d_1(s)} \Theta_s = \Theta'_{m(s)} \circ \alpha_{d_0(s)}$$

because  $\mathcal{C}$  has no 2-cells.



Note: this can be extended to  $\mathcal{C}$  a dbl. cat<sup>dy</sup>

$\Pi_0 : \mathbb{Fam}(C)_0 \rightarrow \underline{\text{Set}}$   
is a split fibration.

We extend this to

$\Pi : \mathbb{Fam}(C) \rightarrow \mathbb{Span}(\underline{\text{Set}})$ .  
(send proarrows to their underlying  
spans and cells to span morphisms)

Claim: this is a split double fibration.

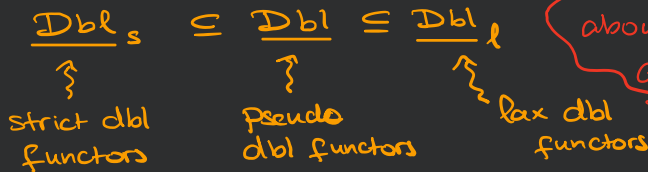


Relation with  
Street's Internal Fibrations  
(a look-off trail we may not take)



# Internal Fibrations

- Internal fibrations in a 2-category were introduced by Street in 1974.
- We will use the following three 2-cat<sup>s</sup>:



say more  
about Street's  
definition?

- In all 3 cases we use vertical transformations

$$\alpha: F \Rightarrow G; \mathbb{D} \Rightarrow \mathbb{E}:$$

\* for each  $x$  in  $\mathbb{D}$ : an arrow  $\alpha_x: Fx \rightarrow Gx$  in  $\mathbb{E}$

\* for each pb arrow  $x \xrightarrow{m} y$  a dbl

cell  $\alpha_m: Fm \rightarrow Gm$  in  $\mathbb{E}$ . (natural and functorial)





## Theorem (Cruttwell, Lambert, P, Szylid)

A strict double functor is an internal fibration in DblCat if and only if it is a double fibration

In addition, a pseudo double functor  $P$

- is an internal fibration in DblCat <sub>$\mathcal{E}$</sub>  iff  $P_0$  and  $P_1$  admit cleavages that are preserved by  $s_{\mathcal{E}}$  and  $t_{\mathcal{E}}$ .
- is an internal fibration in DblCat iff, in addition,  $y_{\mathcal{E}}$  and  $\otimes_{\mathcal{E}}$  preserve Cartesian morphisms.



Furthermore, a strict double functor  $\mathcal{P}$  is an internal fibration in  $\underline{\text{DblCat}}_s$  if and only if  $\mathcal{P}_0$  and  $\mathcal{P}_1$  are fibrations that admit cleavages that are preserved by all of  $s_E$ ,  $t_E$ ,  $y_E$  and  $\otimes_E$



# Double Indexing Functors

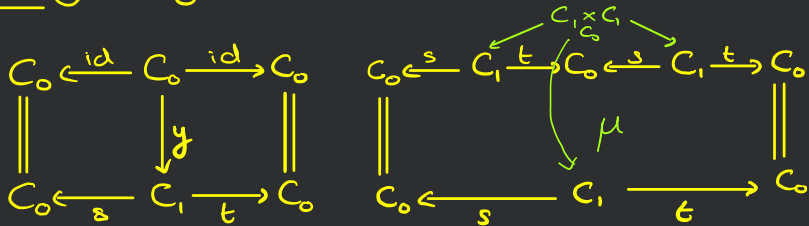
(Take 1)



# Double Indexing Functors

## Exercise!

Note: (1) Categories are monoids in  $\text{Span}(\text{set})$



Unit  $\leftrightarrow$  identities  
for the  $\text{cat}^y$

multiplication  $\leftrightarrow$   
composition for the  
 $\text{cat}^y$ .

(2) Moeller and Vasilakopoulou used:

$$\underline{\text{Fib}} \simeq \underline{\text{ICat}} \Rightarrow \underline{\text{Ps Mon}}(\underline{\text{Fib}}, x) \simeq \underline{\text{Ps Mon}}(\underline{\text{ICat}}, x)$$



To generalize this further we need:

\* double 2-categories ( pseudo category objects in 2-Cat )

\* pseudo monoids in double 2-categories

Result: pseudo categories in a 2-cat<sup>y</sup>  $\mathcal{C}$   
correspond to pseudo monoids in  $\text{Span}(\mathcal{C})$ .

Recall that we want to take the source and target from a more restricted class of arrows, say  $\Sigma$ :

Result: pseudo cat<sup>s</sup> in  $\mathcal{C}$  with  $s, t$  in  $\Sigma$  correspond to pseudo monoids in  $\text{Span}_{\Sigma}(\mathcal{C})$ .

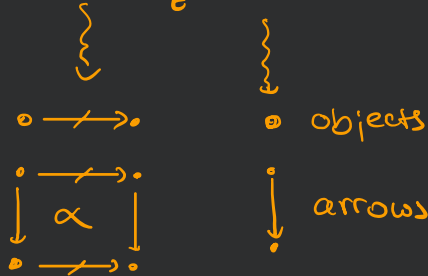


# Double 2-Categories

A double 2-category is a pseudo cat<sup>d</sup> in 2Cat

$$\mathcal{E}_1 \times_{\mathcal{E}_0} \mathcal{E}_1 \xrightarrow{\otimes} \mathcal{E}_1 \begin{array}{c} \xrightarrow{s} \\ \xleftarrow{t} \\ \xrightarrow{e} \end{array} \mathcal{E}_0$$

so we have



$$\Gamma: \alpha \Rightarrow \beta \quad \left( \begin{array}{c} \alpha \\ \Gamma \\ \beta \end{array} \right) \quad \left( \begin{array}{c} \alpha \\ \Gamma \\ \beta \end{array} \right) \quad \text{2-cells}$$



# The Double 2-Category $\text{Span}(\underline{\text{Cat}})$

Objects categories  $\mathcal{A}, \mathcal{B}, \mathcal{C}, \dots$

Arrows functors

Proarrows spans  $\mathcal{B} \xleftarrow{S} \mathcal{A} \xrightarrow{T} \mathcal{C}$

Dbl cells commutative diagrams

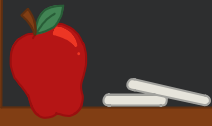
$$\begin{array}{ccccc} \mathcal{B} & \xleftarrow{S} & \mathcal{A} & \xrightarrow{T} & \mathcal{C} \\ \mathcal{G} \downarrow & & \downarrow F & & \downarrow H \\ \mathcal{B}' & \xleftarrow{S'} & \mathcal{A}' & \xrightarrow{T'} & \mathcal{C}' \end{array}$$



# Dbl 3-cells

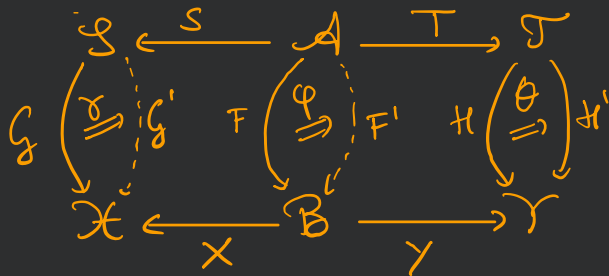
$$\begin{array}{ccccc} \mathcal{S} & \xleftarrow{S} & \mathcal{A} & \xrightarrow{T} & \mathcal{T} & (\gamma, \varphi, \theta) \\ \mathcal{G} \downarrow & & \downarrow F & & \downarrow H & \Rightarrow \\ \mathcal{X} & \xleftarrow{x} & \mathcal{B} & \xrightarrow{y} & \mathcal{Y} & \end{array}$$

$$\begin{array}{ccccc} \mathcal{S} & \xleftarrow{S} & \mathcal{A} & \xrightarrow{T} & \mathcal{T} \\ \mathcal{G}' \downarrow & & \downarrow F' & & \downarrow H' \\ \mathcal{X} & \xleftarrow{x} & \mathcal{B} & \xrightarrow{y} & \mathcal{Y} \end{array}$$





Consists of



such that the two cylinders commute.



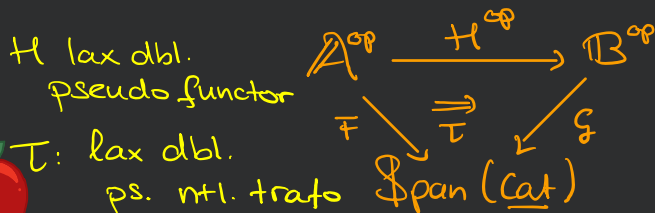
# Double Indexing Functors

$\mathbf{I} \text{Span}(\underline{\text{Cat}})$  is the slice category

with

- objects: Contravariant lax double pseudo functors  $F: \mathcal{A}^{\text{op}} \rightarrow \text{Span}(\underline{\text{Cat}})$  ( $\mathcal{A}$  - indexed ps. dbl. categories)

- morphisms:  $(H, \tau): F \rightarrow G:$



# Lax Double Pseudo Functors

Definition A lax double pseudo functor

$$F : \mathbb{D} \rightarrow \mathbb{E}$$

between dbl 2-categories  $\mathbb{D}, \mathbb{E}$ , consists of:

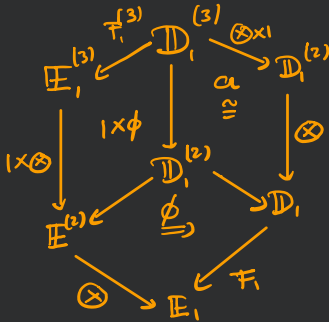
- Pseudo functors :  $F_0 : \mathbb{D}_0 \rightarrow \mathbb{E}_0$  (now 2-categories!)  
 $F_1 : \mathbb{D}_1 \rightarrow \mathbb{E}_1$  pseudo in the arrow direction!
- Comparison pseudo natural transformations:

$$\begin{array}{ccc}
 \mathbb{D}_0 \times \mathbb{D}_0 & \xrightarrow{\otimes} & \mathbb{D}_0 \\
 \downarrow F_0 & \Downarrow \phi & \downarrow F_0 \\
 \mathbb{E}_0 \times \mathbb{E}_0 & \xrightarrow{\otimes} & \mathbb{E}_0
 \end{array}
 \qquad
 \begin{array}{ccc}
 \mathbb{D}_1 & \xrightarrow{\gamma} & \mathbb{D}_1 \\
 \downarrow F_1 & \Downarrow \eta & \downarrow F_1 \\
 \mathbb{E}_1 & \xrightarrow{\gamma} & \mathbb{E}_1
 \end{array}$$

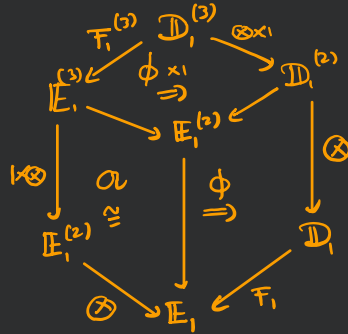
lax in the proarrow direction



- Invertible associativity and unitor modifications



$\cong$



etc.

satisfying well-definedness and coherence conditions.



# Lax Double Pseudo Natural Transformations

A lax dbl ps. ntl transformation:

$$\mathcal{T} : \mathcal{F} \Rightarrow \mathcal{G} : \mathbb{D} \rightrightarrows \mathbb{E}$$

consists of:

- pseudo ntl transformations

$$\mathcal{T}_0 : \mathcal{F}_0 \Rightarrow \mathcal{G}_0, \quad \mathcal{T}_1 : \mathcal{F}_1 \Rightarrow \mathcal{G}_1$$

(arrow-component transformations)

- modifications:

$$\begin{array}{ccc}
 \mathbb{D}_1 \times_{\mathbb{D}_0} \mathbb{D}_1 & \xrightarrow{\otimes} & \mathbb{D}_1 \\
 \downarrow \mathcal{F}_0 \times \mathcal{F}_1 & \xrightarrow{\mathcal{T}_0 \times \mathcal{T}_1} & \downarrow \mathcal{G}_1 \\
 \mathbb{E}_1 \times_{\mathbb{E}_0} \mathbb{E}_1 & \xrightarrow{\otimes} & \mathbb{E}_1
 \end{array}
 \quad \mathcal{T} \quad
 \begin{array}{ccc}
 \mathbb{D}_1 \times_{\mathbb{D}_0} \mathbb{D}_1 & \xrightarrow{\otimes} & \mathbb{D}_1 \\
 \downarrow \mathcal{F}_0 \times \mathcal{F}_1 & \xrightarrow{\cong} & \downarrow \mathcal{F}_1 \\
 \mathbb{E}_1 \times_{\mathbb{E}_0} \mathbb{E}_1 & \xrightarrow{\otimes} & \mathbb{E}_1
 \end{array}$$



and

$$\begin{array}{ccc}
 \mathbb{D}_0 \xrightarrow{\gamma} \mathbb{D}_1 & & \mathbb{D}_0 \xrightarrow{\gamma} \mathbb{D}_1 = \mathbb{D}_1 \\
 \downarrow F_0 \quad \tau_0 \downarrow \cong \downarrow G_0 & \cong & \downarrow F_0 \quad \tau_0 \downarrow \cong \downarrow G_0 \\
 \mathbb{E}_0 \xrightarrow{\beta} \mathbb{E}_1 & & \mathbb{E}_0 \xrightarrow{\beta} \mathbb{E}_1 = \mathbb{E}_1
 \end{array}$$

satisfying multiplicativity and  
 unitality conditions.

We write  $\underline{\text{Dbl2Cat}}(\mathbb{D}, \mathbb{E})$  for the  $\text{cat}^{\gamma}$   
 of lax dbl pseudo functors and lax dbl. ps. ntl.  
 transformations.



# The Representation Theorem



# The Representation Theorem

Theorem (Cruttwell, Lambert, P., Szjld)

There is an equivalence of categories

$$\underline{\text{DblFib}} \simeq \mathbf{I} \text{Span}(\underline{\text{Cat}})$$

Idea for the proof: use pseudo monoids in double 2-categories.

$$\underline{\text{Fib}} \simeq \underline{\text{ICat}} \rightsquigarrow \underline{\text{CFib}} \simeq \underline{\text{ICat}}_t$$

$$\text{so: } \text{Span}_c(\underline{\text{Fib}}) \simeq \text{Span}_t(\underline{\text{ICat}})$$

(convince yourself!)





$\text{Span}_c(\underline{\text{Fib}})$  has objects:  $\begin{array}{c} E \\ \downarrow p \\ B \end{array}$  fibrations with cleavage Notation:  $\begin{array}{c} P \\ \downarrow f \\ P' \end{array}$

$\text{Span}_f(\underline{\text{Cat}})$  has objects:  $\underline{B}^{\text{op}} \xrightarrow{F} \underline{\text{Cat}}$  ps. functor  $F$

arrows:  $\begin{array}{ccc} E & \xrightarrow{f^T} & E' \\ p \downarrow & & \downarrow p' \\ B & \xrightarrow{f^+} & B' \end{array}$  Cartesian arrow preserving  $\begin{array}{c} P \\ \downarrow f \\ P' \end{array}$

arrows:  $\begin{array}{ccc} \underline{B}^{\text{op}} & \xrightarrow{H} & (\underline{B}')^{\text{op}} \\ \downarrow F & \xRightarrow{\theta} & \downarrow F' \\ & \underline{\text{Cat}} & \end{array}$   $\theta$  pseudo trafo  $(H, \theta)$   $\begin{array}{c} F \\ \downarrow \\ F' \end{array}$



Proarrows:  $P \xleftarrow{\ell} Q \xrightarrow{r} R$   $\ell^T, r^T$  cleavage preserving

Proarrows:  $F \xleftarrow{(L, \lambda)} G \xrightarrow{(R, \rho)} K$   $\lambda, \rho$  strict natural transf<sup>ns</sup>

cells: 
$$\begin{array}{ccccc} P & \xleftarrow{\ell} & Q & \xrightarrow{r} & R \\ f \downarrow & & \downarrow g & & \downarrow h \\ P' & \xleftarrow{\ell'} & Q' & \xrightarrow{r'} & R' \end{array}$$
  $\ell^T, r^T, \ell'^T, r'^T$   
cleavage preserving  
 $f^T, g^T, h^T$

cells: 
$$\begin{array}{ccccc} F & \xleftarrow{(L, \lambda)} & G & \xrightarrow{(R, \rho)} & K \\ (H, \theta) \downarrow & & \downarrow & & \downarrow \\ F' & \xleftarrow{(L', \lambda')} & G' & \xrightarrow{(R', \rho')} & K' \end{array}$$
 Cartesian arrow pres.



Then lift :

$$\begin{aligned}\underline{\text{Dbf Fib}} &:= \underline{\text{Ps Mon}} (\text{span}_c(\underline{\text{Fib}})) \simeq \\ &\simeq \underline{\text{Ps Mon}} (\text{span}_f(\underline{\text{Icat}})) \simeq \underline{\text{Ispan}}(\underline{\text{Cat}}).\end{aligned}$$



# Connections with Known Constructions

Monoidal Fibrations

monoids  
are cats

Double Fibrations

2-Fibration

quintet constr.

Discrete Dbl Fib.

Double Gr. Constr.

All of these have an indexed notion and a Grothendieck category of elements.

In our case the composition of the functors used in the proof gives us also a category of elements construction.



The Double Category  
of Elements



## The Double Grothendieck Construction

Start with  $F: \mathbb{D}^{\text{op}} \longrightarrow \text{Span}(\underline{\text{Cat}})$ :

$$F_0: \mathbb{D}_0^{\text{op}} \longrightarrow \text{Span}(\underline{\text{Cat}})_0 = \underline{\text{Cat}} \quad (1)$$

$$F_1: \mathbb{D}_1^{\text{op}} \longrightarrow \text{Span}(\underline{\text{Cat}})_1$$

and a further induced functor:

$$\mathbb{D}_1^{\text{op}} \xrightarrow{F_1} \text{Span}(\underline{\text{Cat}})_1 \xrightarrow{\text{apx}} \underline{\text{Cat}} \quad (2)$$

Apply the ordinary elements construction to (1) and (2):

$$\mathbb{E}(F)_0 \longrightarrow \mathbb{D}_0 \quad \mathbb{E}(F)_1 \longrightarrow \mathbb{D}_1$$

cloven fibrations.



## Details of $\mathbb{E}l(\mathbb{F})$

Some notation:

- for  $m: A \rightarrow B$  in  $\mathbb{D}$ , write  $\mathbb{F}m$  in  $\text{Span}(\text{Cat})$   
as:

$$\begin{array}{ccc} & \mathbb{F}m & \\ Lm \swarrow & & \searrow Rm \\ \mathbb{F}A & & \mathbb{F}B \end{array}$$

- for  $A \xrightarrow{m} B \xrightarrow{n} C$  in  $\mathbb{D}$ , we have the laxity  
comparison cell

$$\begin{array}{ccccc} \mathbb{F}A & \xleftarrow{\quad} & \mathbb{F}m \times_{\mathbb{F}B} \mathbb{F}n & \xrightarrow{\quad} & \mathbb{F}C \\ \parallel & & \downarrow \phi_{m,n} & & \parallel \\ \mathbb{F}A & \xleftarrow{L_{m \circ n}} & \mathbb{F}(m \circ n) & \xrightarrow{R_{m \circ n}} & \mathbb{F}C \end{array}$$



Now  $\mathbb{E}l(F)$  is given by:

• objects:  $(C, x)$   $C$  in  $\mathbb{D}$ ,  $x$  in  $F_C$ .

• arrows:  $(f, \bar{f}): (C, x) \rightarrow (D, y)$

with  $f: C \rightarrow D$  in  $\mathbb{D}$

and  $\bar{f}: x \rightarrow f^*y (= F(f)y)$  in  $F_C$ .

• proarrows:  $(m, \bar{m}): (C, x) \rightrightarrows (D, y)$

with  $C \xrightarrow{m} D$  in  $\mathbb{D}$

$$\bar{m} \in F_m \quad \left( F_C \xleftarrow{L_m} F_m \xrightarrow{R_m} F_D \right)$$

$$\text{s.t. } L_m(\bar{m}) = x$$

$$R_m(\bar{m}) = y$$





• double cells:

$$\begin{array}{ccc}
 (A, x) & \xrightarrow{(m, \bar{m})} & (B, y) \\
 (s, \bar{f}) \downarrow & (\theta, \bar{\theta}) & \downarrow (g, \bar{g}) \\
 (C, z) & \xrightarrow{(n, \bar{n})} & (D, w)
 \end{array}$$

with

$$\begin{array}{ccc}
 A & \xrightarrow{m} & B \\
 f \downarrow & \theta & \downarrow g \\
 C & \xrightarrow{n} & D
 \end{array}$$

$$\begin{array}{ccccc}
 FA & \xleftarrow{L_m} & Fm & \xrightarrow{R_m} & FB \\
 F_f \downarrow & & \downarrow F_\theta & & \downarrow F_g \\
 FC & \xleftarrow{L_n} & Fn & \xrightarrow{R_n} & FD
 \end{array}$$

in  $\mathcal{D}$

and  $\bar{m} \xrightarrow{\bar{\theta}} \theta^* \bar{n}$  an arrow in  $Fm$  s.t.

$$L_m(\bar{\theta}) = \bar{f} \quad \text{and} \quad R_m(\bar{\theta}) = \bar{g}.$$



Composition in the "arrow direction" is as expected:

- for arrows  $(A, x) \xrightarrow{(f, \bar{f})} (B, y) \xrightarrow{(g, \bar{g})} (C, z)$   
the composite is  $(gf, \phi_{f, g} f^*(\bar{g}) \bar{f}) : (A, x) \rightarrow (C, z)$

- for cells  $(m, \bar{m}) \xRightarrow{(\theta, \bar{\theta})} (n, \bar{n}) \xRightarrow{(\delta, \bar{\delta})} (p, \bar{p})$

the composite is

$$(\delta\theta, \phi_{\theta, \delta} \theta^*(\bar{\delta}) \bar{\theta}) : (m, \bar{m}) \Rightarrow (p, \bar{p}).$$

- Units:  $(1_C, (\varphi_C)_x) : (C, x) \rightarrow (C, x)$   
 $(1_m, (\varphi_m)_{\bar{m}}) : (m, \bar{m}) \rightarrow (m, \bar{m}).$



$\mathbb{E}l(\mathcal{F}) \begin{array}{c} \xrightarrow{s} \\ \xrightarrow{t} \end{array} \mathbb{E}l(\mathcal{F})_0$  are defined by

$$s(\theta, \bar{\theta}) = (f, \bar{f})$$

$$t(\theta, \bar{\theta}) = (g, \bar{g})$$

- proarrow composition:

$$\text{for } (A, x) \xrightarrow{(m, \bar{m})} (B, y) \xrightarrow{(n, \bar{n})} (C, z)$$

the composite is:

$$(m \otimes n, \phi_{m,n}(\bar{m}, \bar{n})) : (A, x) \longrightarrow (C, z)$$

- Composition for cells and proarrow units are given using appropriate components of the structure maps related to pseudo naturality.



## Results

- $\Pi: \mathbb{E}l(\mathcal{F}) \rightarrow \mathbb{D}$  is a double fibration.
- This is the object part of an equivalence of categories

$$\underline{\text{Dbl Fib}} \simeq \mathbf{I} \text{Span}(\underline{\text{Cat}})$$

which specializes to

$$\underline{\text{Dbl Fib}}(\mathbb{B}) \simeq \underline{\text{Dbl 2Cat}}(\mathbb{B}^{\text{op}}, \text{Span}(\underline{\text{Cat}}))$$

for each dbl cat<sup>d</sup>  $\mathbb{B}$ .



## Examples

- Let  $\underline{A}$  be a category with pushouts and  $\mathcal{Csp}(\underline{A})$  the double cat<sup>y</sup> with cells:

$$\begin{array}{ccccc} X & \longrightarrow & Z & \longleftarrow & Y \\ u \downarrow & & \downarrow w & & \downarrow v \\ X' & \longrightarrow & Z' & \longleftarrow & Y' \end{array}$$

- For any lax double functor

$$F: \mathcal{Csp}(\underline{A}) \rightarrow \mathcal{Span}(\underline{Cat})$$

the dbl cat<sup>y</sup> of elements  $|\mathcal{EL}(F)|$  is  $F\text{-Csp}$ ,  
the double category of  $F$ -decorated cospans.  
[Patterson, 2023].



This slightly generalizes the decorated cospans  
from Baez-Courser - Vasilakopoulou.