

MARKOV CATEGORIES

A TUTORIAL

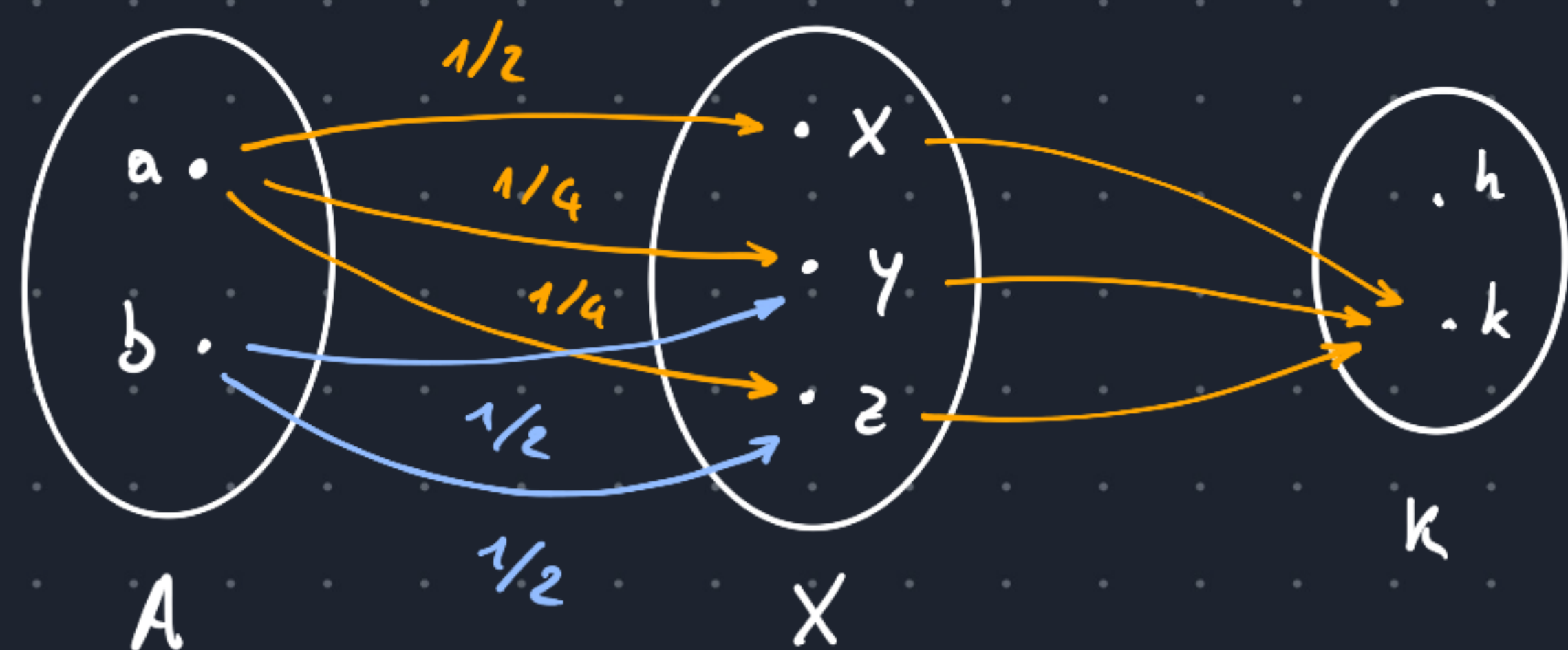
Applied Category Theory 2023

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Basic idea

Morphisms with "randomness"!

Example. The category Fin Stoch whose morphisms are stochastic matrices.



	a	b
x	1/2	0
y	1/4	1/2
z	1/4	1/2

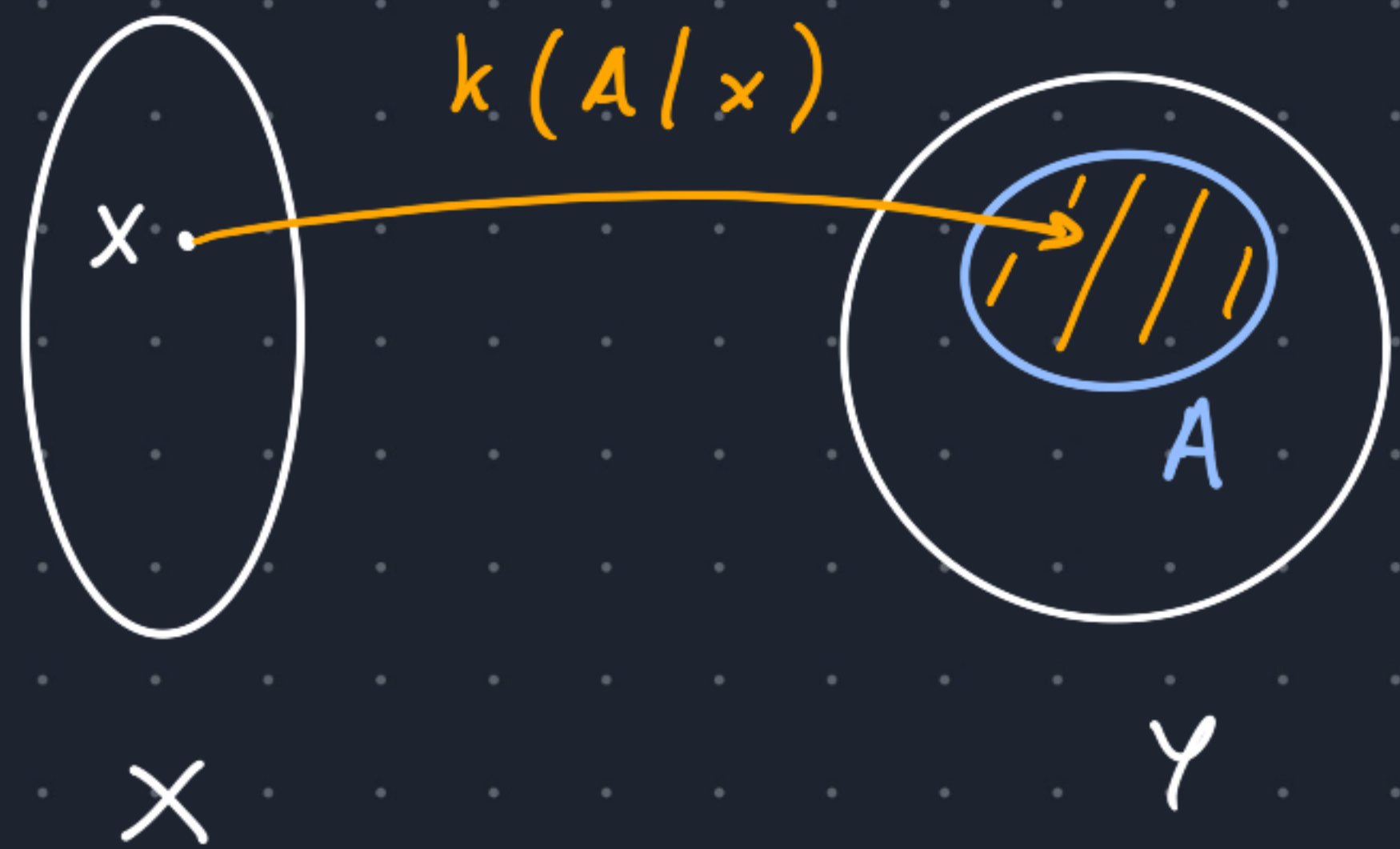
$f(z|b)$

$$g \circ f(k|a) = \sum_{x \in X} g(k|x) f(x|a) \quad (\text{Chapman-Kolmogorov})$$

$$1 \xrightarrow{P} X$$

"states" = prob. measures (of finite support)

Example. The category *Stoch* whose morphisms are Markov kernels.



$$X \times \Sigma_Y \longrightarrow [0,1]$$

$$(x, A) \longmapsto k(A|x)$$

prob. measure
measurable

$$h \circ k(B|x) = \int_Y h(B|y) k(dy|x)$$

$$X \xrightarrow{f} Y \quad K_f(A|x) := \begin{cases} 1 & f(x) \in A \\ 0 & f(x) \notin A \end{cases} \quad \text{Meas} \xrightarrow{K} \text{Stoch}$$

Main definition

A Markov category is a symmetric monoidal category $(\mathcal{C}, \otimes, I)$, where each object X is equipped with maps

$X \xrightarrow{\text{copy}} X \otimes X$

$X \xrightarrow{\text{del}} I$



such that



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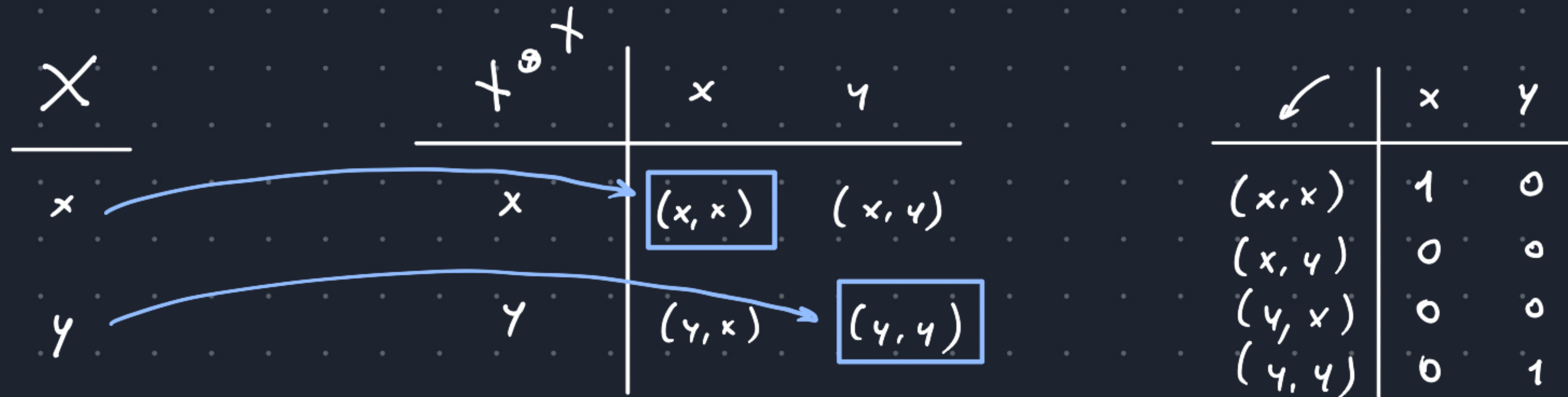


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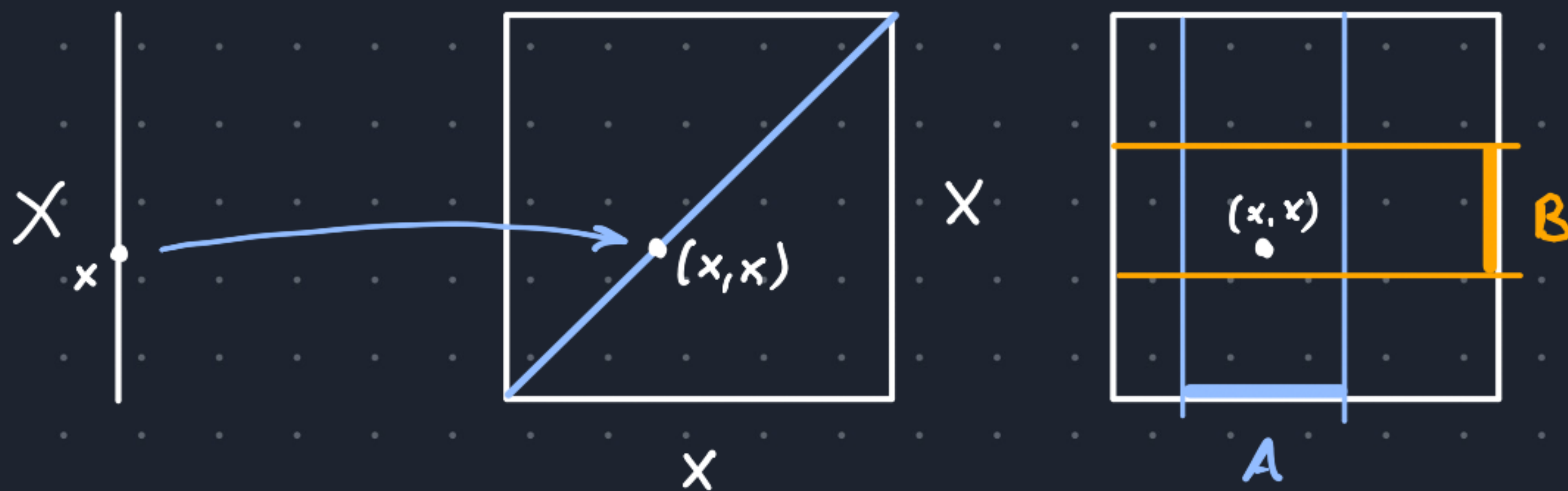


Without this:
"GS" or "CD"
category

Example. In FinStoch, $X \xrightarrow{\text{copy}} X \otimes X$ for $X = \{x, y\}$ is:



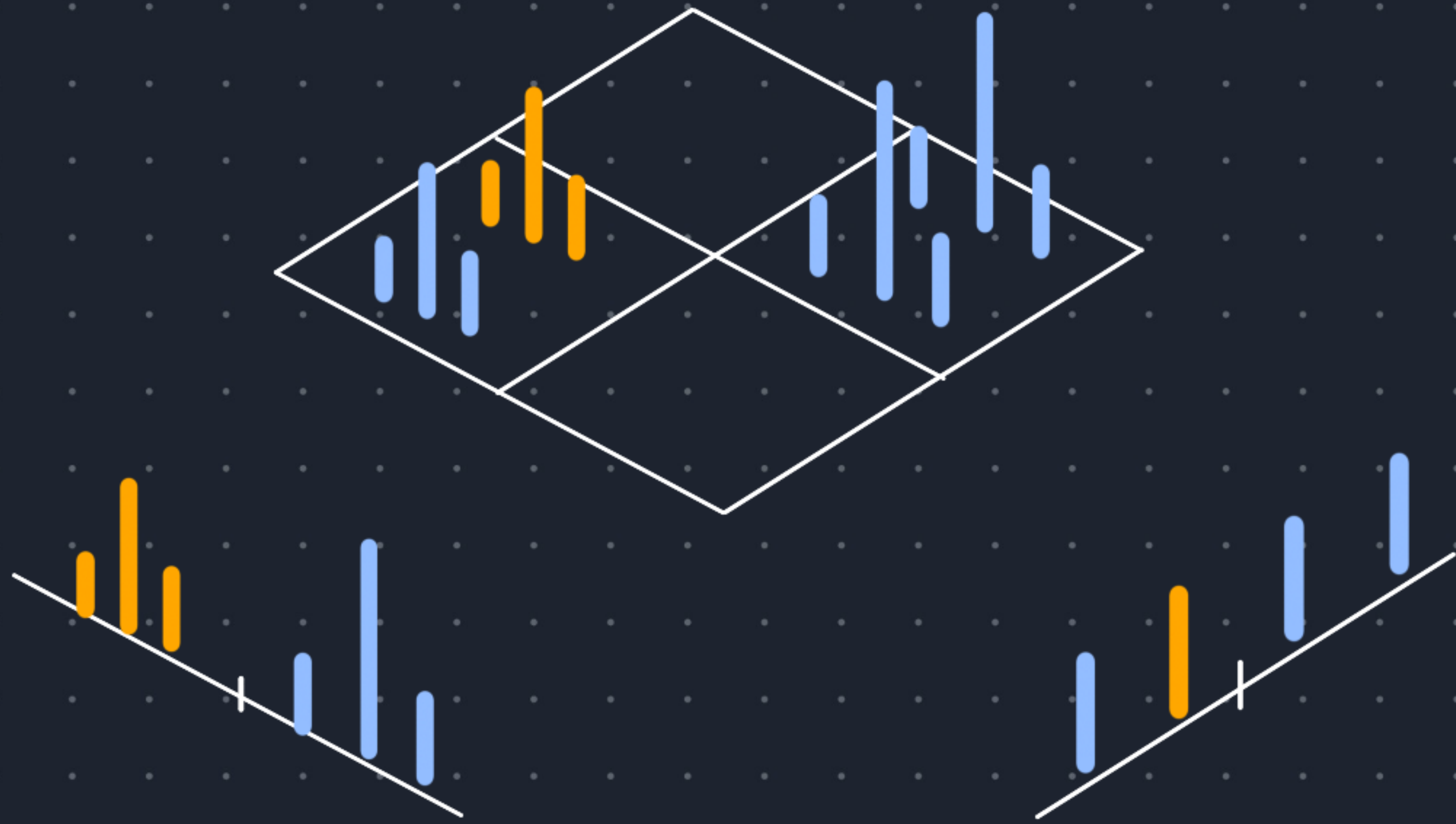
Example. In Stoch, more generally,



$$(x, x) \in A \times B \iff x \in A \cap B$$

$$\text{copy}(A \times B | x) = \begin{cases} 1 & x \in A \cap B \\ 0 & x \notin A \cap B \end{cases}$$

Joints & Marginals



$$\sum_y r(x, y) = p(x)$$

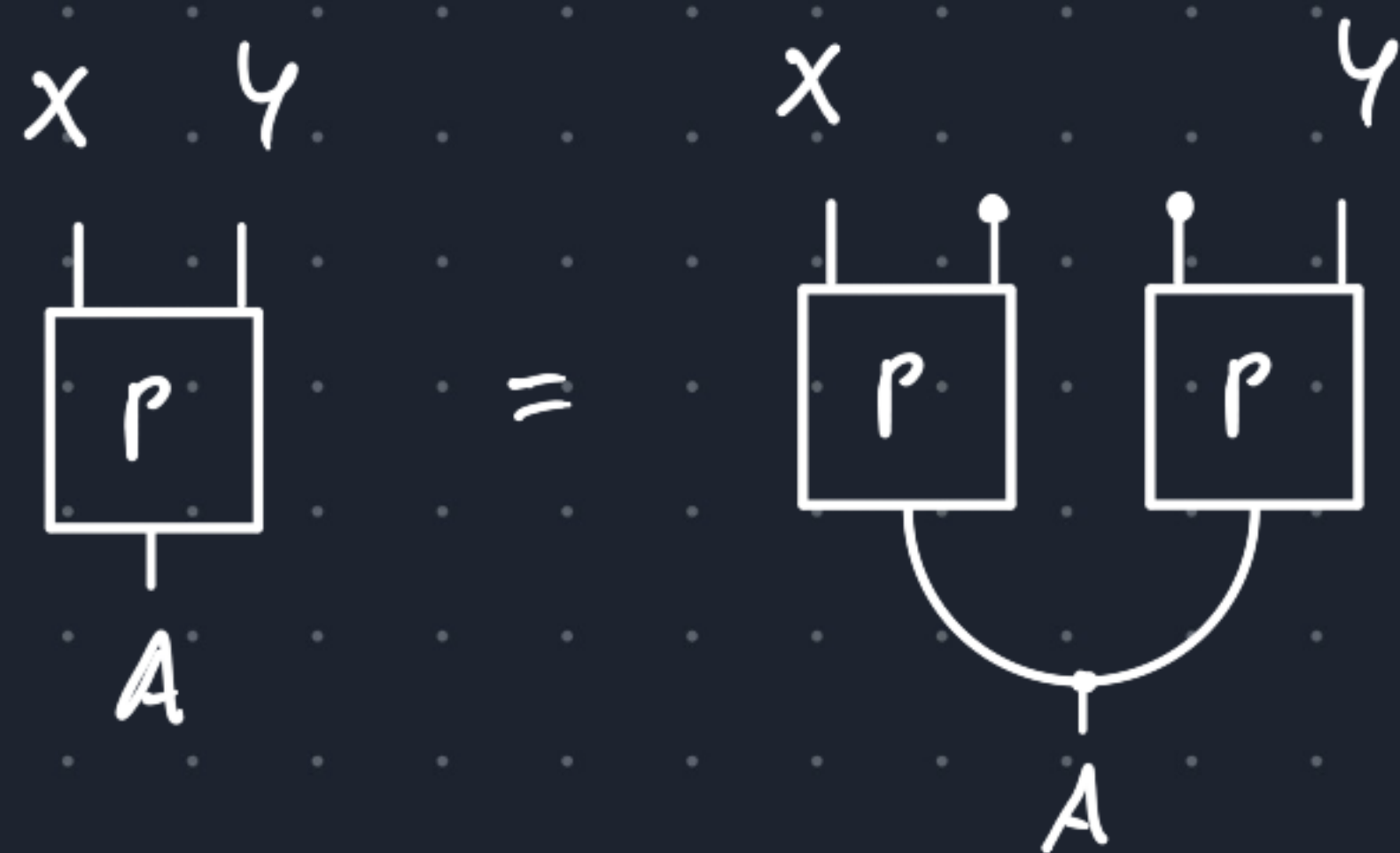
(Same for y)

Stochastic independence



p exhibits independence of x, y

$$p(x, y) = p(x) p(y)$$



h exhibits conditional independence
of x, y given A .

$$p(x, y | a) = p(x | a) p(y | a)$$

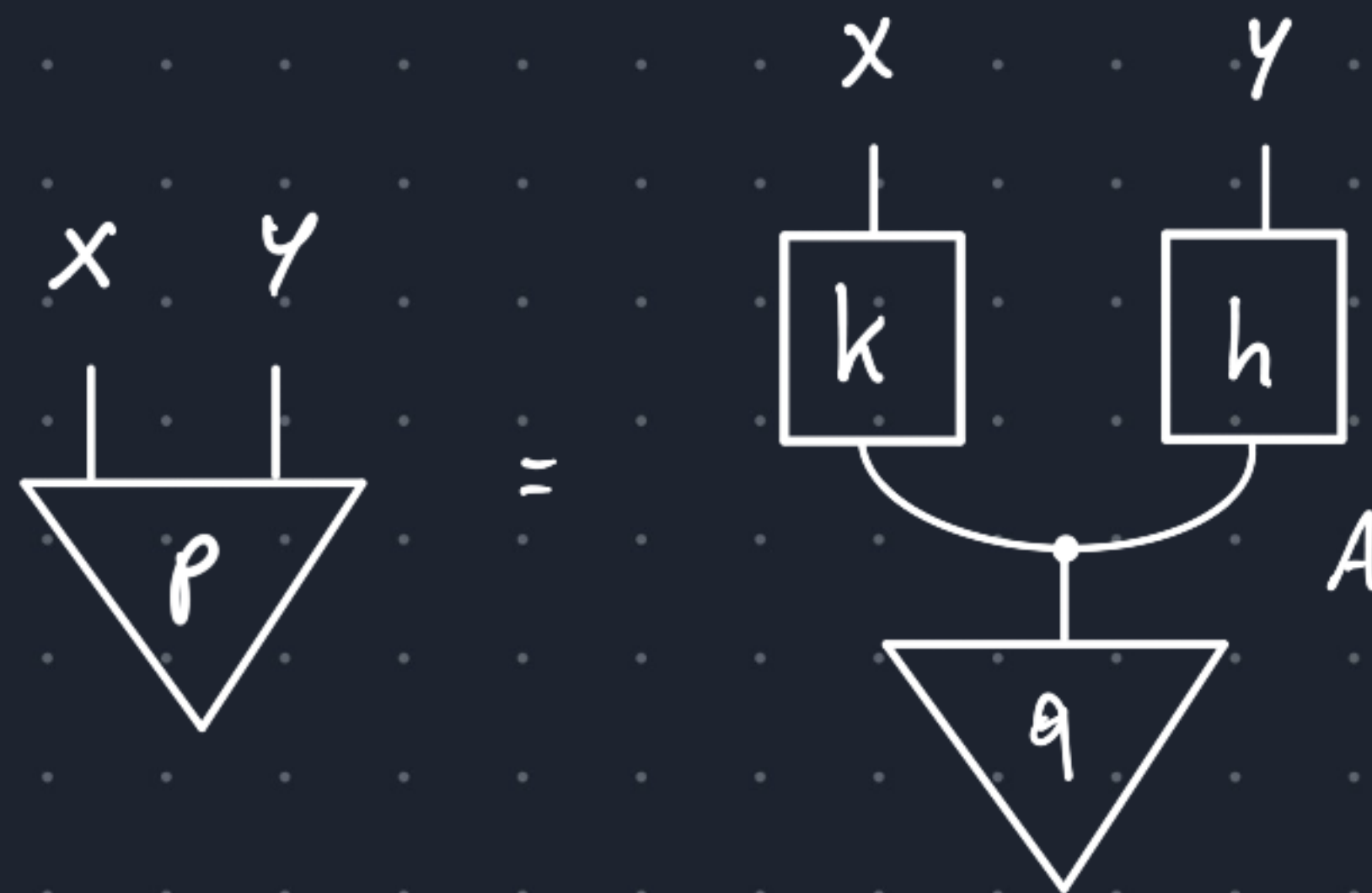
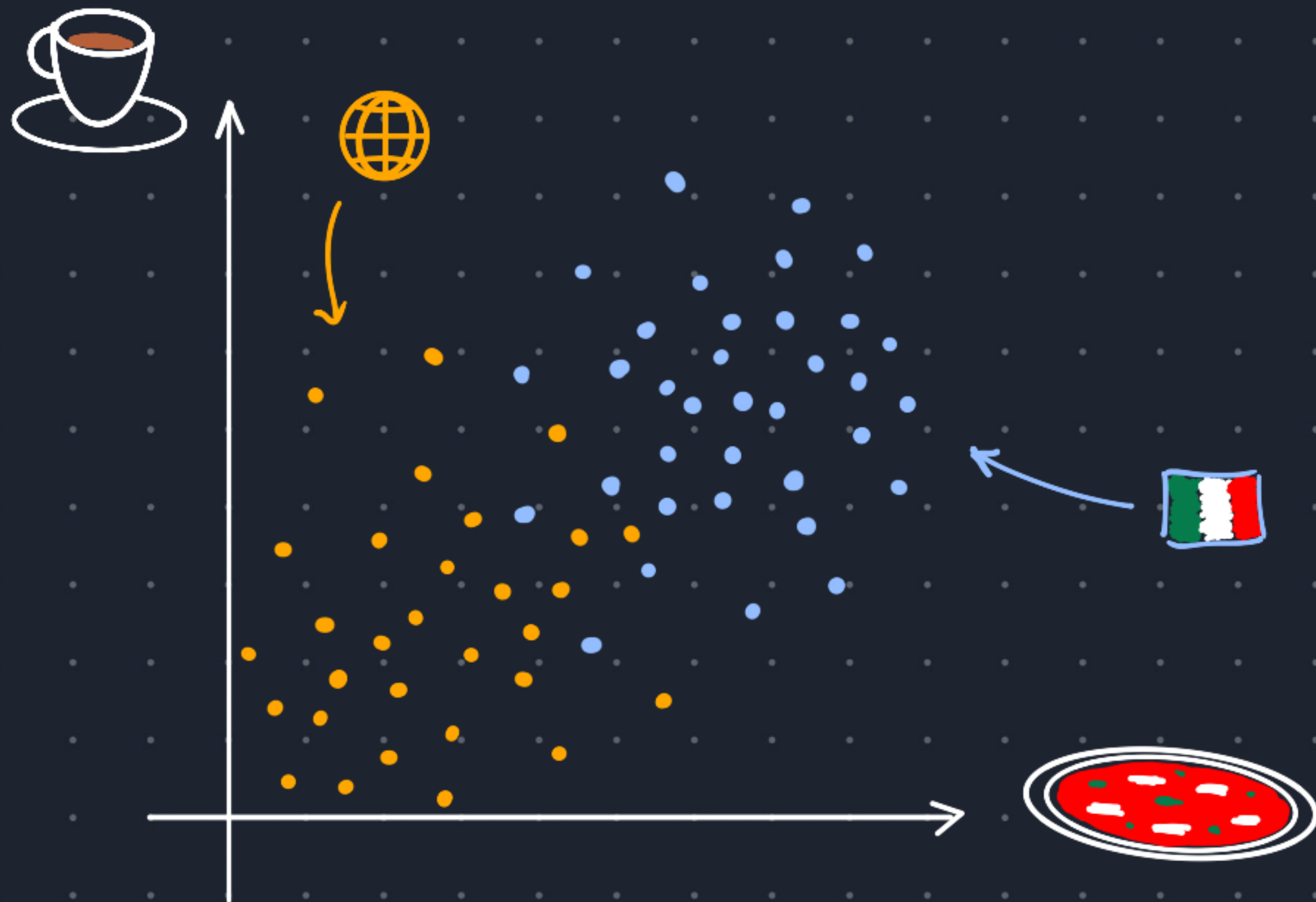
(See Exercise 1.1 later for more.)

Stochastic independence



p exhibits independence of x, y

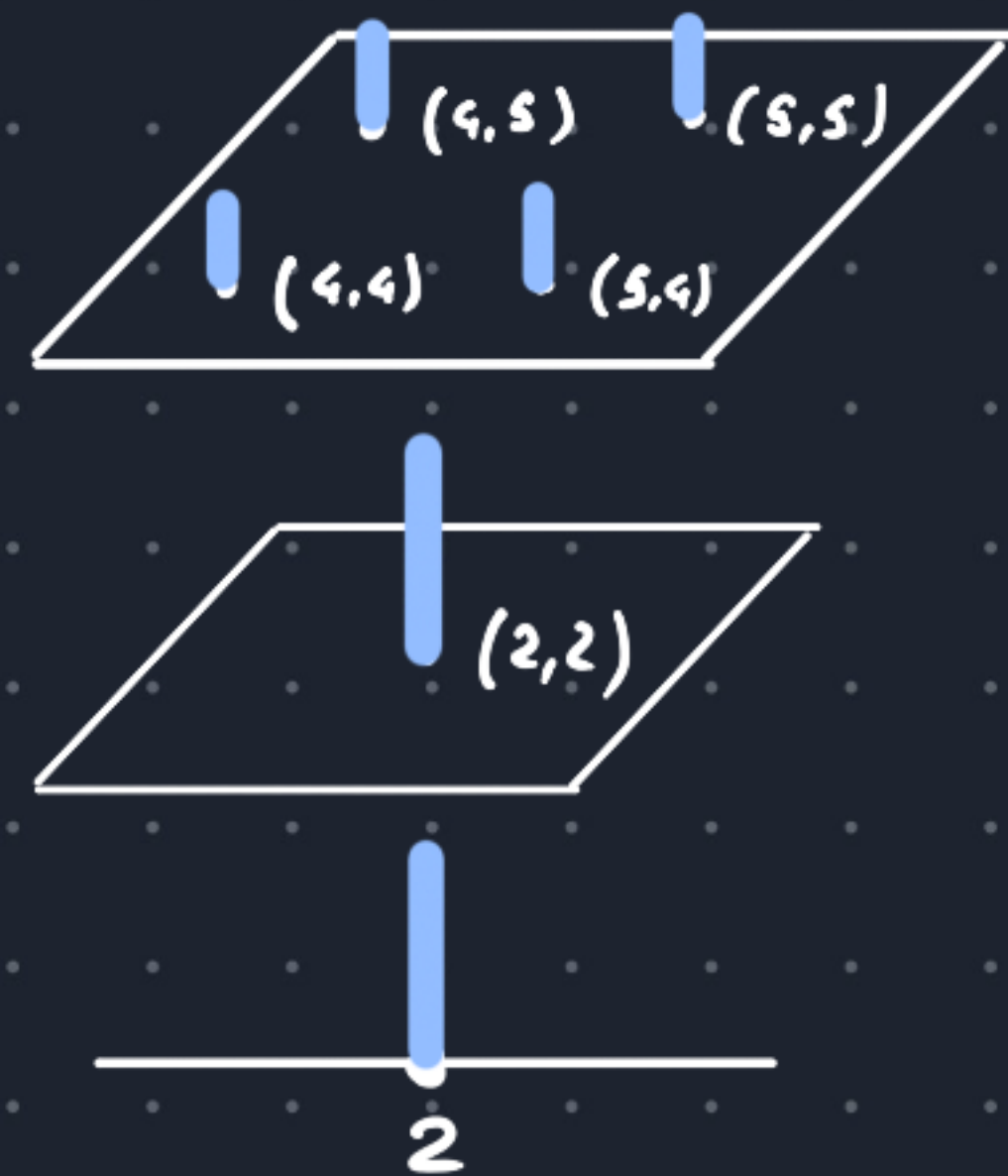
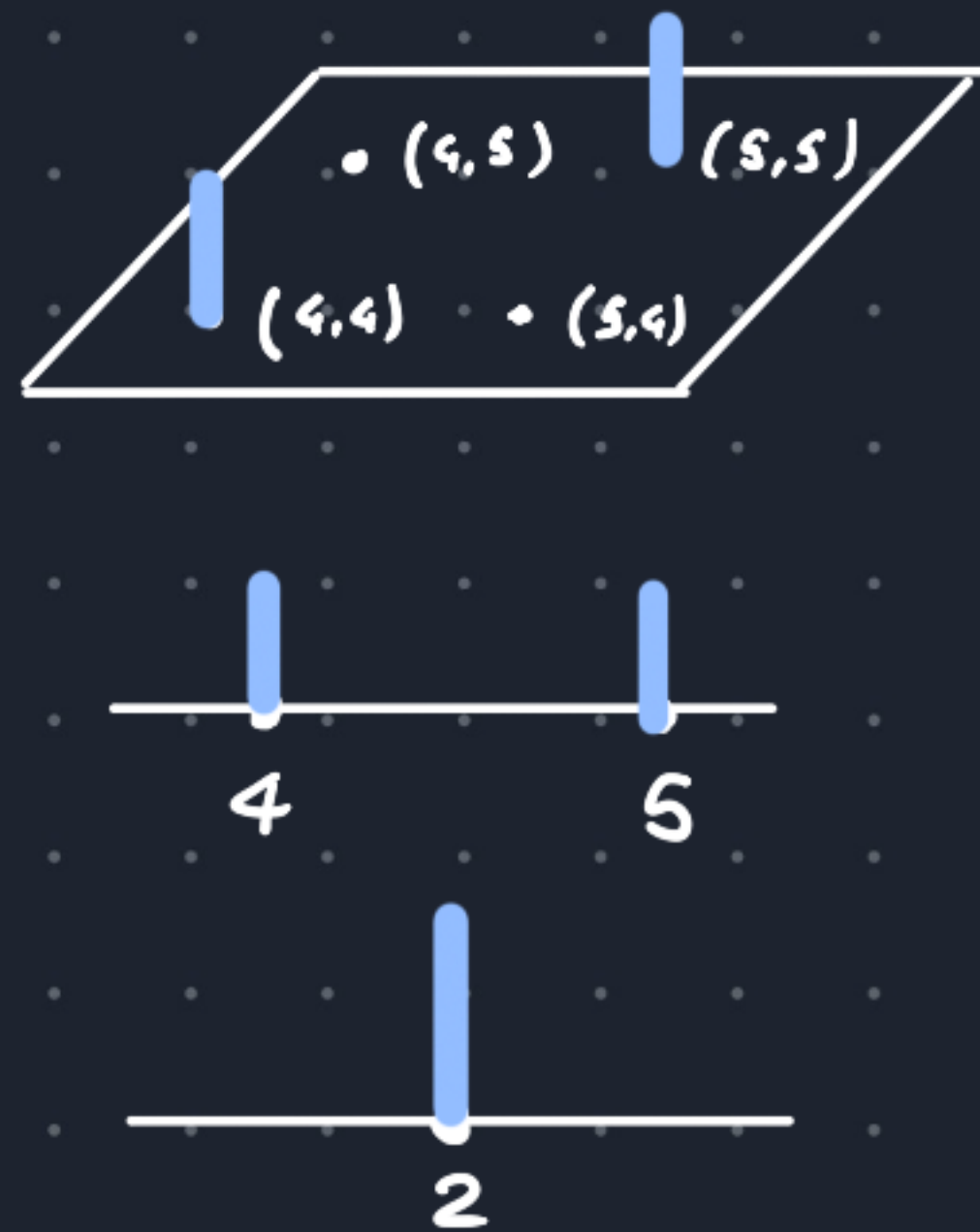
$$p(x, y) = p(x) p(y)$$



Determinism (a.k.a. copyability)



is called deterministic if



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is called deterministic if



=



Example.

- In FinStoch, a matrix is deterministic iff all its entries are $f(y|x) = 0$ or 1 .
- In Stoch, similarly, $f(A|x) = 0$ or 1 .

(See Exercise 1.2 later.)



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Determinism (a.k.a. copyability)



is called deterministic if



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- The category **BorelStoch** is the subcategory of standard Borel spaces.
(e.g. finite & countable sets, $\mathbb{R} \cong [0,1] \cong \mathbb{R}^n$, etc.)

In this category, deterministic morphisms are just the measurable functions.

$$K_f(A|x) := \begin{cases} 1 & f(x) \in A \\ 0 & f(x) \notin A \end{cases}$$

Determinism (a.k.a. copyability)

Proposition. For a Markov category \mathcal{C} , TFAE:

- 1) Every morphism is deterministic;
- 2) The copy maps are a natural transformation;
- 3) \mathcal{C} is cartesian monoidal ($\otimes = \times$)

Markov = cartesian + randomness!
cartesian = Markov + determinism



Stochastic interaction is
a feature of randomness.

Exercises:

1.1. Show that if a joint morphism decomposes

as a product, i.e.

The diagram shows a large box labeled 'h' with two input wires from below labeled 'A'. Two output wires go upwards, labeled 'X' and 'Y'. This is shown to be equal to two smaller boxes labeled 'f' and 'g' side-by-side. Each has one input wire from below labeled 'A'. The output of 'f' is wire 'X' and the output of 'g' is wire 'Y'. A curved line connects the two 'A' inputs to the two boxes.

then it is the product of its marginals, i.e.

The diagram shows a large box labeled 'h' with two input wires from below labeled 'A'. Two output wires go upwards, labeled 'X' and 'Y'. This is shown to be equal to two smaller boxes labeled 'h' side-by-side. Each has one input wire from below labeled 'A'. The output of the left 'h' is wire 'X' and the output of the right 'h' is wire 'Y'. A curved line connects the two 'A' inputs to the two boxes.

1.2. Show that the deterministic morphisms of FinStoch are exactly the matrices of entries $\{0, 1\}$.

What's the analogous statement in Stoch?

1.3. Prove

Proposition. For a Markov category \mathcal{C} , TFAE:

- 1) Every morphism is deterministic;
- 2) The copy maps are a natural transformation;
- 3) \mathcal{C} is cartesian monoidal ($\otimes = \times$)

Hint: show that if a morphism

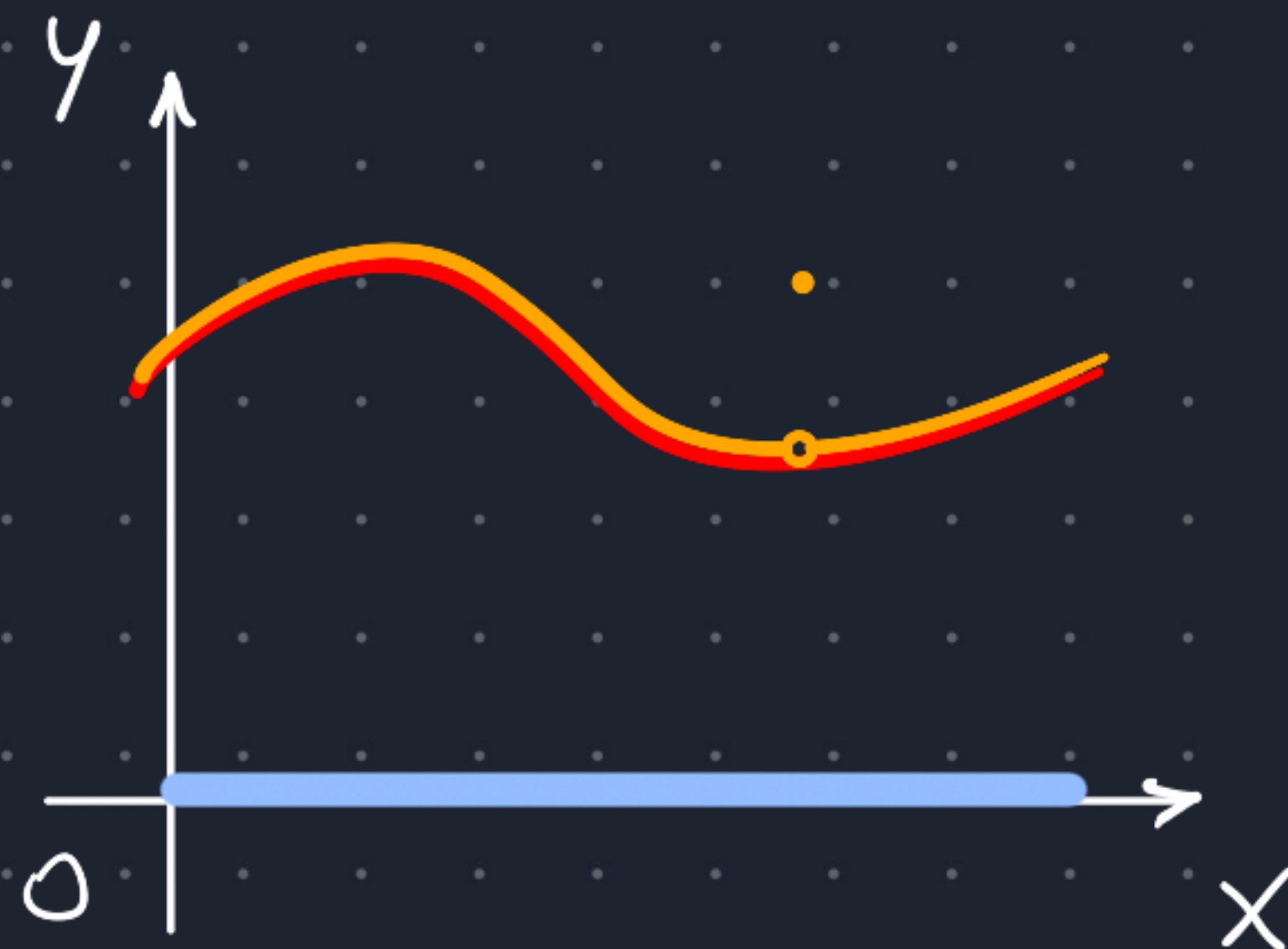
The diagram shows a box labeled 'h' with one input wire from below labeled 'A'. Two output wires go upwards, labeled 'X' and 'Y'.

is deterministic, then it is always making X and Y conditionally independent given A .

Almost-sure equality

Given $\begin{array}{c} x \\ | \\ \triangle \\ P \end{array}$ and $\begin{array}{c} y \\ | \\ \square \\ f \\ | \\ x \end{array}, \begin{array}{c} y \\ | \\ \square \\ g \\ | \\ x \end{array}$, we say that $f = g$ p -almost surely

if $\begin{array}{c} x \quad y \\ | \quad | \\ \cup \quad \square \\ | \quad f \\ | \quad | \\ \triangle \\ P \end{array} = \begin{array}{c} x \quad y \\ | \quad | \\ \cup \quad \square \\ | \quad g \\ | \quad | \\ \triangle \\ P \end{array}$



Example. In FinStoch, $f(x) = g(x)$ at each x s.t. $p(x) \neq 0$.

In BorelStoch, the set $\{x \in X ; f(x) \neq g(x)\}$ has p -measure zero.

(See Exercise 2.1 later.)

Conditioning & Bayesian inverses

Given , a conditional distribution of p given x

is a morphism $*$  such that



=



(* see Exercise 2.2 later)

$$p(x, y) = p(x) p(y|x)$$

Theorem. FinStoch has all conditional distributions.

BorelStoch too.  Traditionally the chosen category for "classical" prob. theory.

Stoch, in general, does not.

Conditioning & Bayesian inverses

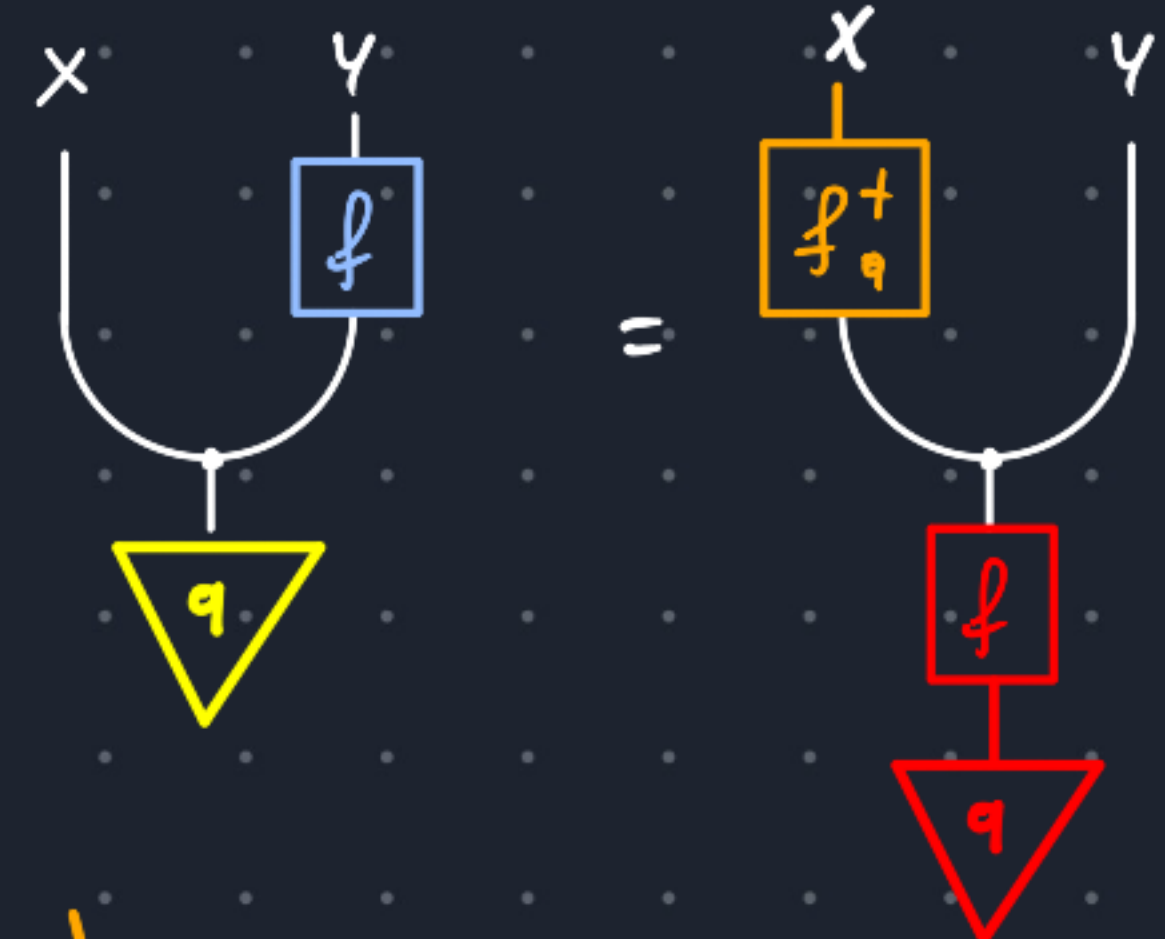
Given , a conditional distribution of p given x

is a morphism  such that



Given  and , a Bayesian inverse of f relative to q , is a conditional dist.

for the joint , ie. a morphism  such that



$$p(x) p(y|x) = p(y) p(x|y)$$

The ProbStoch construction

"Cat. of probability spaces & transport plans"

Definition.

Let \mathcal{C} be a Markov category with all conditional distributions.

The category $\text{ProbStoch}(\mathcal{C})$ has:

• As objects, pairs $(X, \begin{array}{c} x \\ | \\ \triangle \\ \text{p} \end{array})$ where $X \in \mathcal{C}$, $p: I \rightarrow X$
(e.g. prob. spaces)

• As morphisms $(X, \begin{array}{c} x \\ | \\ \triangle \\ \text{p} \end{array}) \rightarrow (\gamma, \begin{array}{c} y \\ | \\ \triangle \\ \text{q} \end{array})$, equivalence classes

under p.a.s. equality of morphisms $\begin{array}{c} \gamma \\ | \\ \square \\ \text{f} \\ | \\ X \end{array}$ such that $\begin{array}{c} \gamma \\ | \\ \square \\ \text{f} \\ | \\ \triangle \\ \text{p} \end{array} = \begin{array}{c} y \\ | \\ \triangle \\ \text{q} \end{array}$.

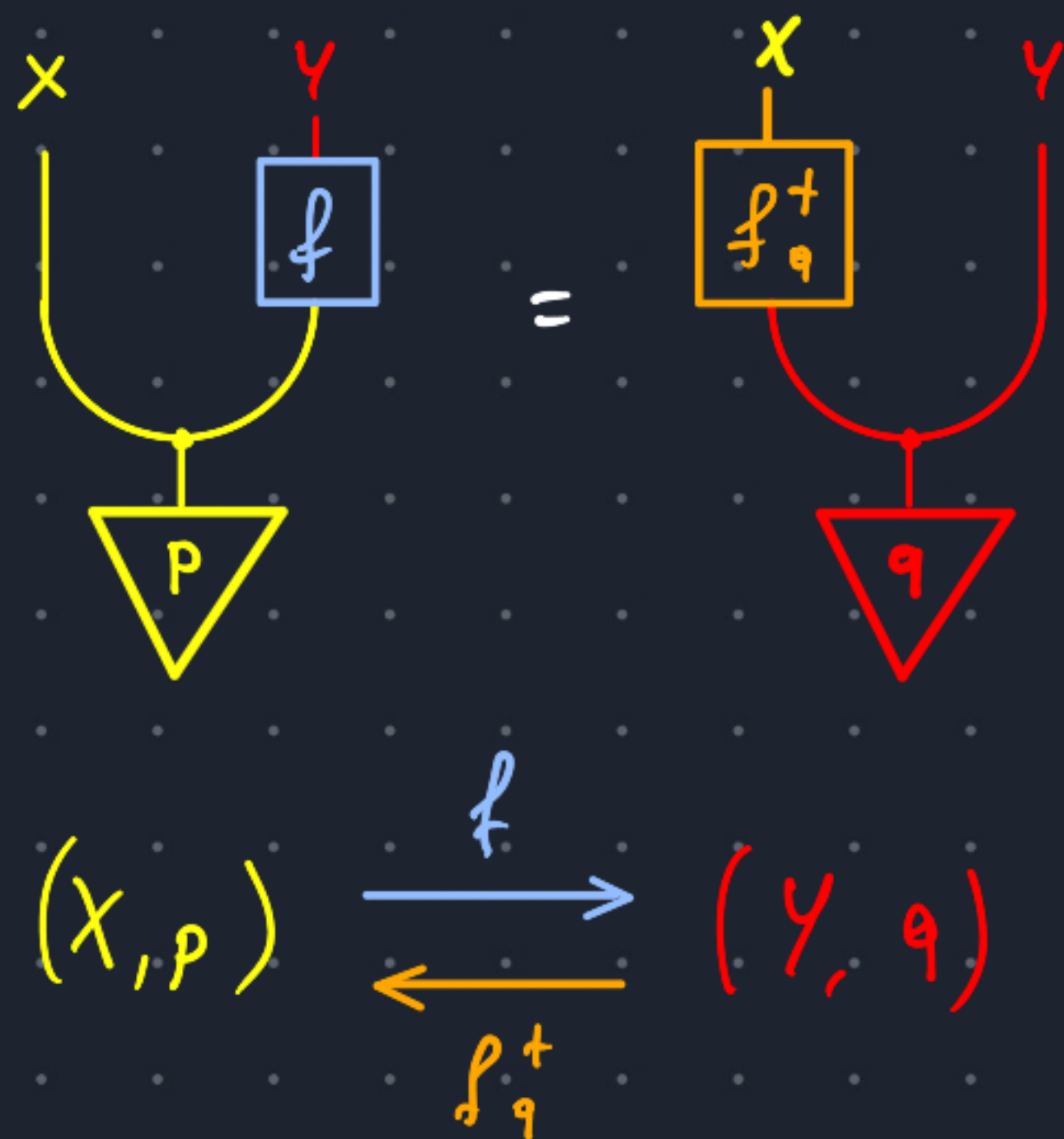
(See Exercise 2.4 later.)

(measure-preserving)

The ProbStoch construction

"Cat. of probability spaces & transport plans"

Theorem. ProbStoch(\mathcal{C}) is a dagger category, where $\dagger =$ Bayesian inversion.



Definition. A dagger structure on a category \mathcal{C} is a "self-duality" functor $\mathcal{C} \xrightarrow{\cong} \mathcal{C}^{\text{op}}$

which is

- Identity on objects: $X^\dagger = X$
- Involutive: $f^{\dagger\dagger} = f$.

Equivalently, morphisms of ProbStoch(\mathcal{C}) are couplings: joint distributions



Exercises:

2.1. Show that, in FinStoch , BorelStoch ,

given $p: I \rightarrow X$ and $f, g: X \rightarrow Y$,

we have that $f = g$ p -almost surely

iff $p(\{x \in X; f(x) \neq g(x)\}) = 0$.

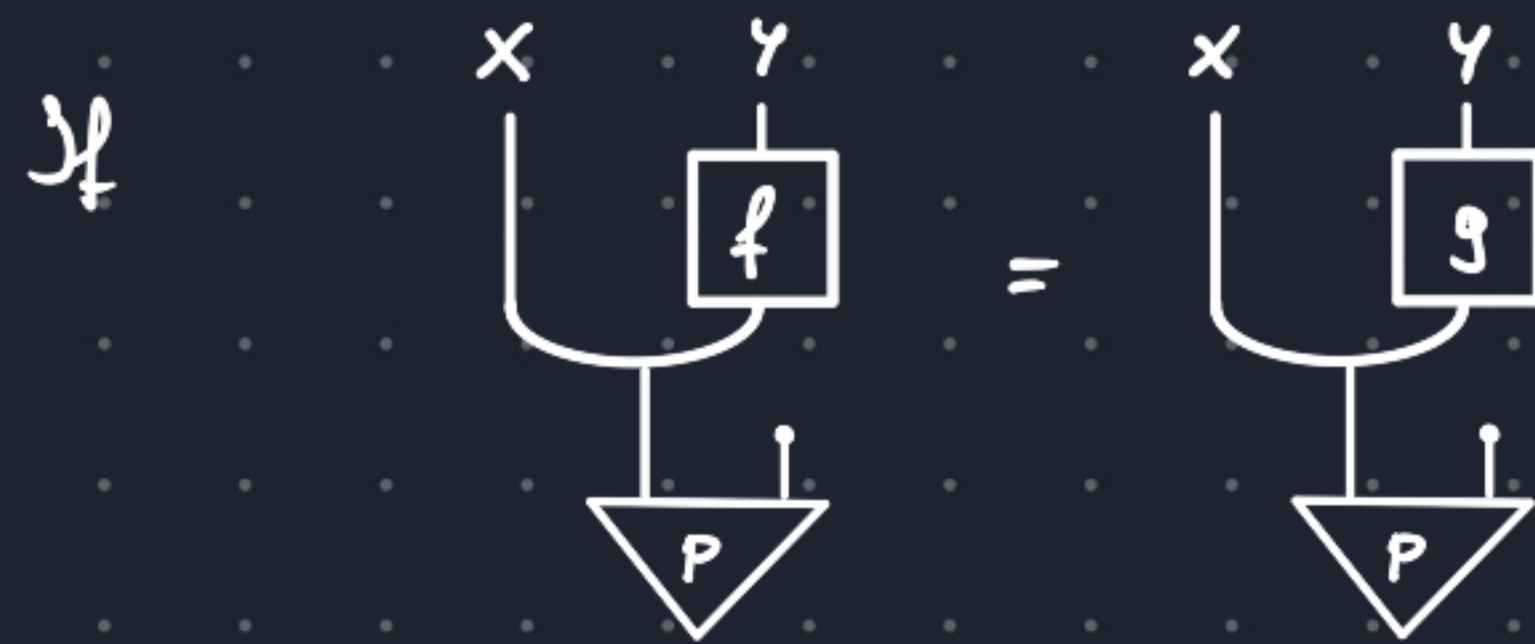
2.2. Suppose the following cond. dist. exists:



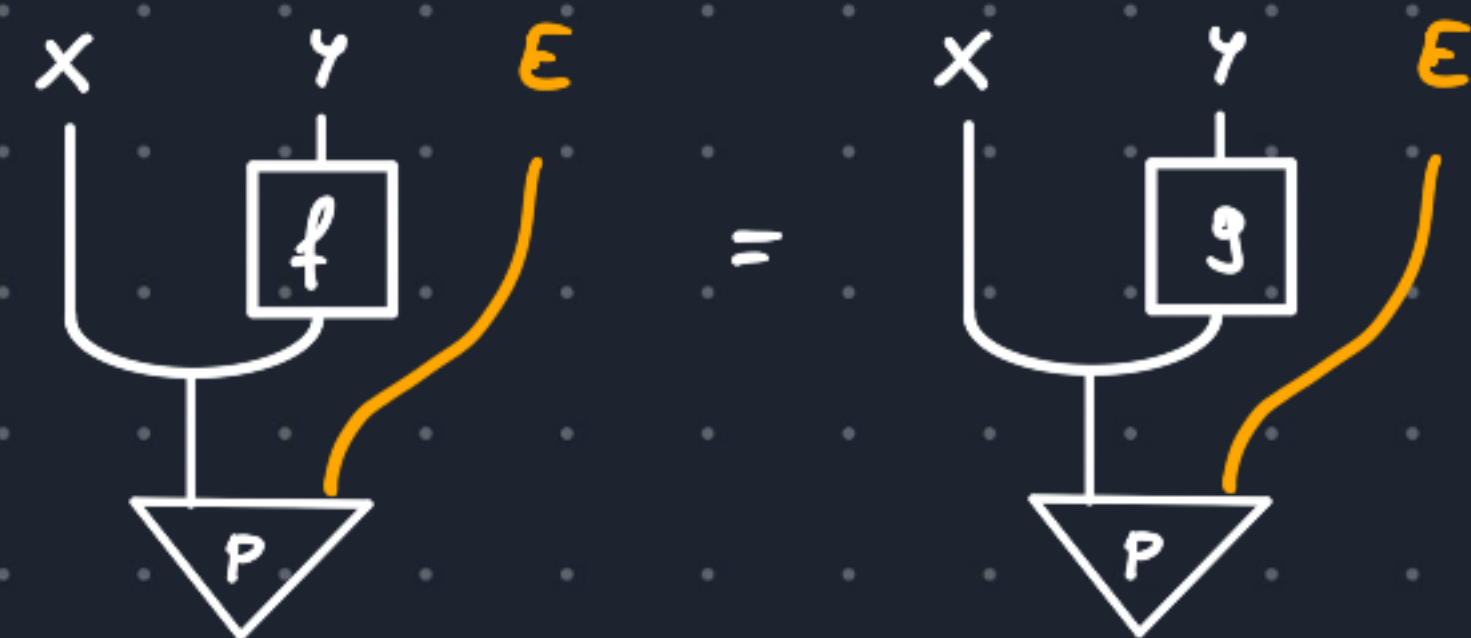
Show that any other morphism $\begin{matrix} Y \\ | \\ f \\ | \\ X \end{matrix}$ is a conditional iff it is p -a.s. equal to $\boxed{p|_x}$.

2.3. Suppose \mathcal{C} has conditional distributions.

Prove the **equality strengthening** property:



then also



2.4. Use Ex. 2.3 to show that composition in $\text{ProbStoch}(\mathcal{C})$ is well defined.

Markov categories and monads

Proposition. Let \mathcal{D} be cartesian monoidal.

Let (P, μ, η) be a monad on \mathcal{D} which is

- Affine : $P1 \cong 1$

- Monoidal (= commutative) : $PA \times PB \xrightarrow{\nabla} P(A \times B)$

Then $\text{Kleisli}(P)$ is a Markov category.

Example. Stoch is the Kleisli category of the Giry monad on Meas . (Exercise 3.2)

FinStoch is almost the Kleisli cat. of the distribution monad on Set .

Markov categories and monads

$$\mathcal{D}(A, PB) \cong \mathcal{D}_p(A, B)$$

$$\mathcal{D}(PB, PB) \cong \mathcal{D}_p(PB, B)$$

$$\text{id} \longmapsto \text{Samp}$$



In basic probability theory:

$$\text{Bernoulli}(p) = \begin{cases} 1 & \text{with prob. } p \\ 0 & \text{with prob. } 1-p \end{cases}$$

$$\underbrace{[0, 1]}_{\text{"2}} \xrightarrow{\text{Bernoulli}} \{0, 1\}$$

$$P(\{0, 1\})$$

The unit of the adjunction is

$$\mathcal{D}(A, PA) \cong \mathcal{D}_p(A, A)$$

$$\delta \longleftarrow \text{id}$$

$$X \xrightarrow{\delta} PX$$



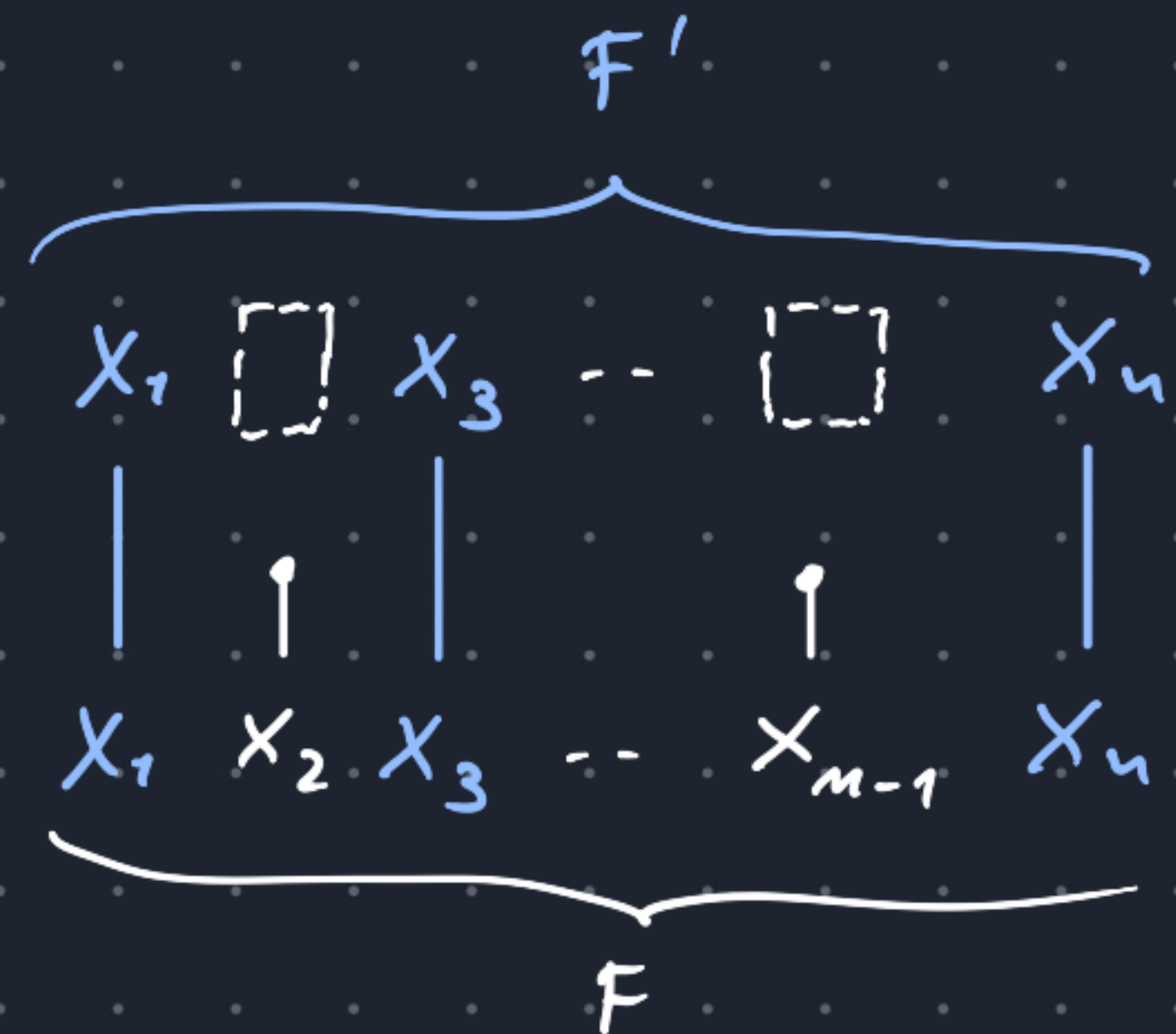
Kolmogorov products

- Probability theory is all about stochastic processes in infinite time, which need objects in the form $X^{\mathbb{N}}$ (at least!).
- While cartesian products can be infinite, we need infinite monoidal products.

• For finite sets F , we have $\bigotimes_{i \in F} X_i$.

• Given subsets $F' \subseteq F$, we can marginalize

$$\bigotimes_{i \in F} X_i \longrightarrow \bigotimes_{i \in F'} X_i \quad \text{by discarding.}$$



Kolmogorov products

Definition. Let \mathcal{C} be a Markov category. Let I be an infinite set, $\{X_i\}_{i \in I}$.

A Kolmogorov product is a cofiltered limit

$$X^I := \lim_{F \subseteq I} \left(\bigotimes_{i \in F} X_i \right) \longrightarrow \dots \begin{array}{ccc} & \nearrow X_i \otimes X_j & \rightarrow X_i \\ & \searrow X_j \otimes X_k & \rightarrow X_j \\ & & \searrow X_k \end{array}$$

such that

• It is preserved by $Y \otimes -$

• The arrows $X^I \longrightarrow X^F$ are deterministic.

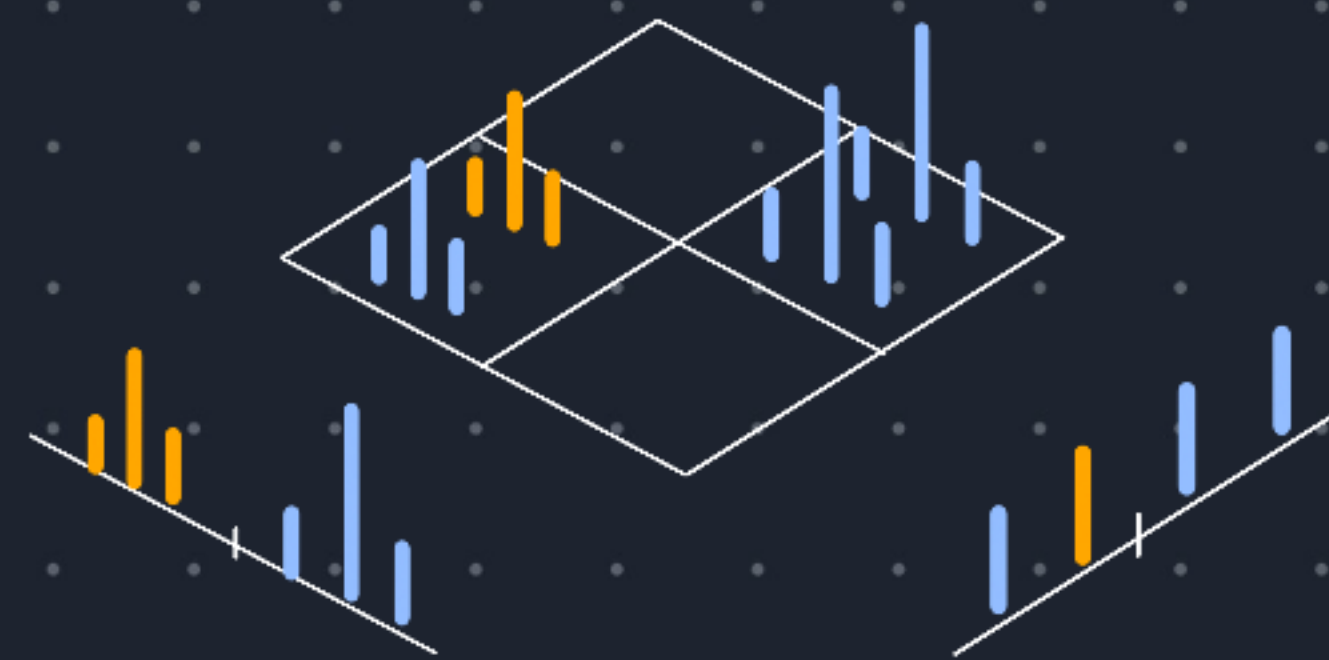
Theorem (Kolmogorov extension). Borel Stoch has countable Kolmogorov products.

Kolmogorov products

When \mathcal{C} is the Kleisli category of a probability monad, a Kolmogorov product encodes the absence of infinitary stochastic interactions.

- "No products": P does not preserve finite products

$$P(X \times Y) \xrightarrow{\neq} P X \times P Y$$



- Kolmogorov ext. thm:

$$P(X^{\mathbb{N}}) = P\left(\lim_{F \subseteq \mathbb{N}} X^F\right) \xrightarrow{\cong} \lim_{F \subseteq \mathbb{N}} P(X^F) \xrightarrow{\neq} \lim_{F \subseteq \mathbb{N}} (P X)^F = (P X)^{\mathbb{N}}$$

Exercises

3.1. Prove

Proposition. Let \mathcal{D} be cartesian monoidal.

Let (P, μ, η) be a monad on \mathcal{D} which is

- Affine : $P1 \cong 1$
- Monoidal : $PA \times PB \xrightarrow{\Delta} P(A \times B)$

Then $\text{Kleisli}(P)$ is a Markov category.

3.2. (For people who know some measure theory.)

Prove that $\text{Stoch} \cong \text{Kleisli}(\text{Giry monad})$

Hint: why is a kernel a Kleisli morphism?

3.3. Prove that sampling from a product distribution is the same as sampling the factors independently:

$$\begin{array}{ccc} PX \otimes PY & \xrightarrow{\Delta} & P(X \otimes Y) \\ & \searrow \text{samp} \otimes \text{samp} & \downarrow \text{samp} \\ & & X \otimes Y \end{array}$$

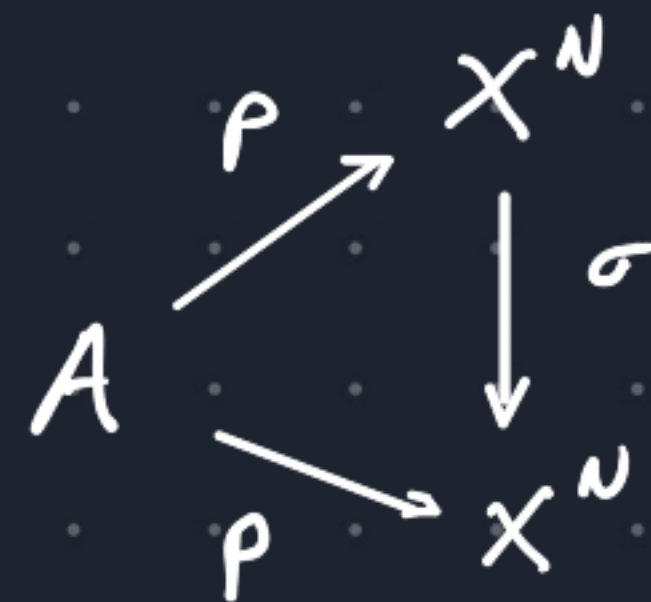
3.4. Prove that the Kolmogorov product

$\bigotimes_{i \in I} X_i$ is the cartesian product

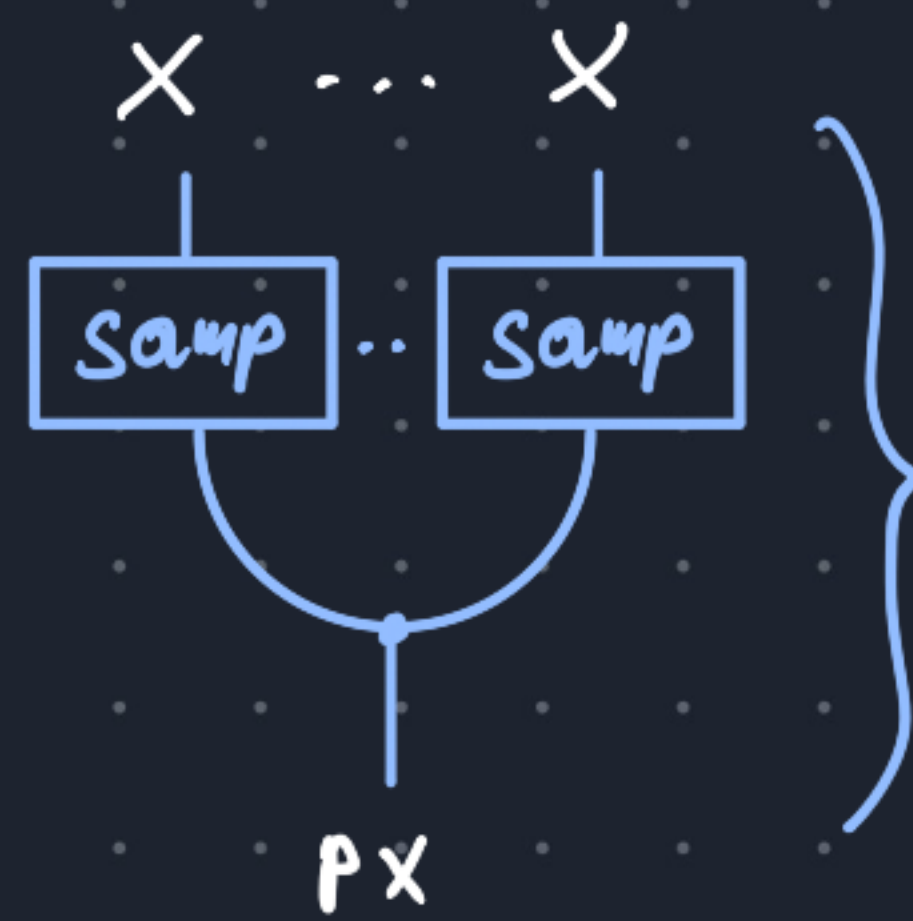
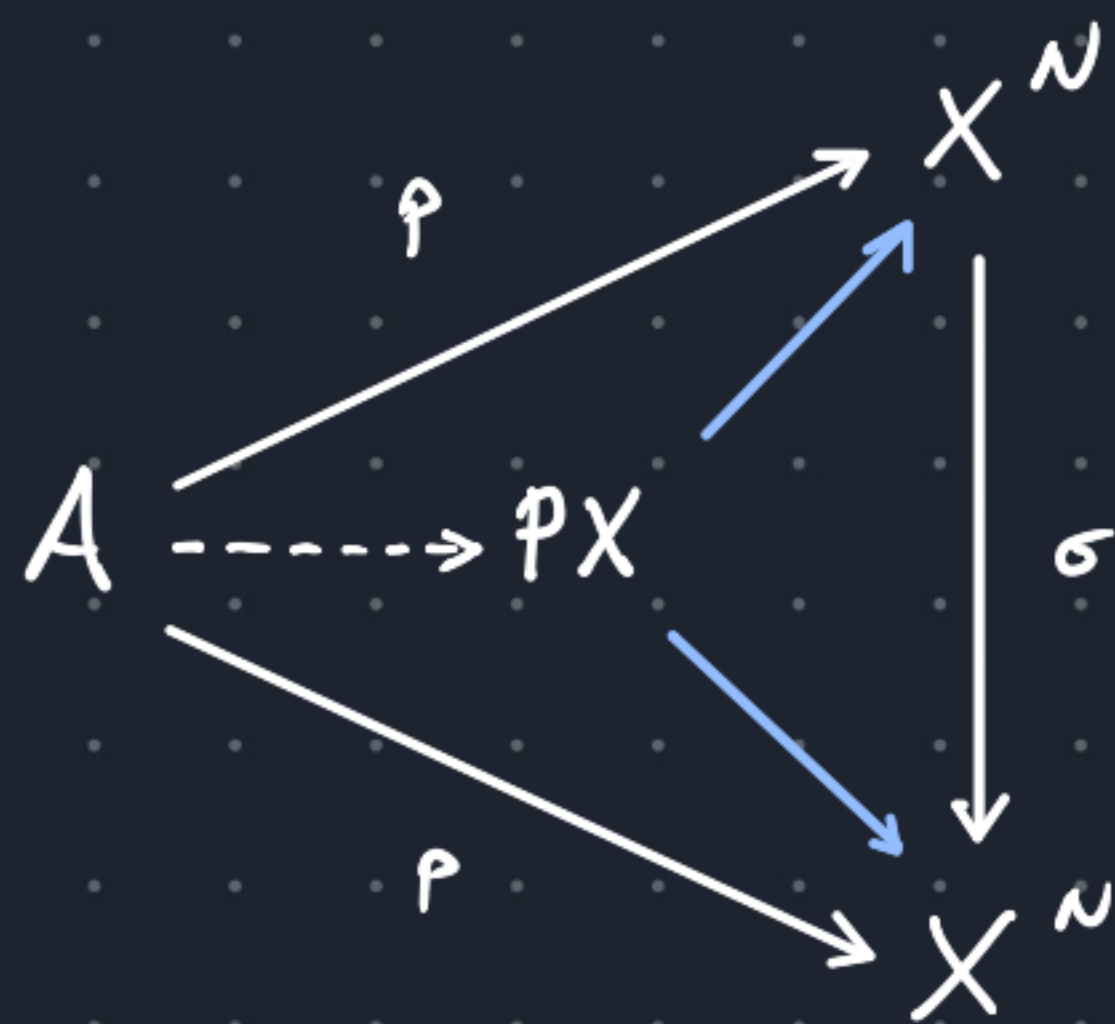
$\prod_{i \in I} X_i$ in the subcategory \mathcal{C}_{def} .

The de Finetti Theorem

A morphism $\begin{array}{c} X^N \\ \square \\ \rho \\ \square \\ A \end{array} = \begin{array}{c} X \dots X \\ \swarrow \quad \searrow \\ \square \\ \rho \\ \square \\ A \end{array}$ is called *exchangeable* if it commutes with finite permutations (in the result).



Theorem. In Borel Stoch, for every A and X , there is a natural bijection



taking i.i.d. samples!
(limiting cone)

The de Finetti Theorem



The X are conditionally independent given P_X (the distribution from which they are sampled - independently).

Example. Let $X = \{\text{Heads}, \text{Tails}\}$. Flip a coin repeatedly (exchangeable).

Suppose you see Heads, Heads, Heads, Heads.

What do you expect to see next?

What if you know that the coin is fair?

"Same coin!"

Results so far

Classical probability:

- De Finetti theorem (Fritz-Gonda-Penone '21)
- d-separation criterion (Fritz-Klingler '22)
- Kolmogorov extension theorem (Fritz-Rischel '19)
- Kolmogorov, H-S O-1 laws (Fritz-Rischel '19)
- Multinomial, hypergeometric distributions (Jacobs '21)

Statistics:

- Theorems on sufficient statistics (Fritz '19)
- Comparison of experiments (Fritz-Gonda-Penone-Rischel '20)

Ergodic theory, information theory:

- Ergodic decomposition thm (Moss-Penone '22)
- Entropy, data processing inequalities (Penone '22)

Theoretical computer science:

- Privacy eqn (Sabok et al '20, Fritz et al. '22)
- Observational monads (Moss-Penone '22)

Quantum probability:

- Quantum Markov categories (Parzygnat '20, '21)

+ more in progress!

Some references:

- K. Cho, B. Jacobs, Disintegration and Bayesian inversion via string diagrams. *Mathematical Structures in Computer Science*. [arXiv: 1709.00322](#)
 - T. Fritz, A synthetic approach to Markov kernels, conditional independence, and theorems on sufficient statistics. *Advances in Mathematics*. [arXiv: 1908.07021](#)
 - T. Fritz, T. Gonda, P. Penone, E. F. Rischel, Representable Markov categories and comparison of statistical experiments in categorical probability. [arXiv: 2010.07416](#)
 - T. Fritz, T. Gonda, P. Penone, The de Finetti theorem in categorical probability. *Journal of Stochastic Analysis*. [arXiv: 2105.02639](#)
 - T. Fritz, T. Gonda, M. Gauguin Houghton-Larsen, P. Penone, D. Stein, Dilations and information flow axioms in categorical probability. [arXiv: 2211.02507](#)
 - P. Penone, Markov Categories and Entropy. [arXiv: 2212.11719](#)
- ... & more